Introduction to Analysis of Algorithms

**Objectives** 

- To introduce the concept of analysing algorithms with respect to the time taken to have them executed
  - Purpose:
    - To see if an algorithm is practical
    - To compare different algorithms for solving a problem

Introduction to Analysis of Algorithms

- One aspect of software quality is the efficient use of *computer resources*:
  - CPU time
  - Memory usage
- We frequently want to analyse algorithms with respect to *execution time* 
  - Called *time complexity* analysis
  - For example, to decide which sorting algorithm will take less time to run

# **Time Complexity**

- Analysis of time taken is based on:
  - Problem size (e.g. number of items to sort)
  - Primitive operations (e.g. comparison of two values)
- What we want to analyse is the relationship between
  - The size of the problem, n
  - And the time it takes to solve the problem, t(n)
    - Note that t(n) is a function of n, so it depends on the size of the problem

# **Time Complexity Functions**

- This t(n) is called a *time complexity function*
- What does a time complexity function look like?
  - Example of a time complexity function for some algorithm:

 $t(n) = 15n^2 + 45 n$ 

 See the next slide to see how t(n) changes as n gets bigger!

#### Example: 15n<sup>2</sup> + 45 n

No. of items <b>n</b>	15n <sup>2</sup>	45n	15n <sup>2</sup> + 45n
1	15	45	60
2	60	90	150
5	375	225	600
10	1,500	450	1,950
100	150,000	4,500	154,500
1,000	15,000,000	45,000	15,045,000
10,000	1,500,000,000	450,000	1,500,450,000
100,000	150,000,000,000	4,500,000	150,004,500,000
1,000,000	15,000,000,000,000	45,000,000	15,000,045,000,000

### Comparison of Terms in 15n<sup>2</sup> + 45 n

- When n is small, which term is larger?
- But, as n gets larger, note that the 15n<sup>2</sup> term grows more quickly than the 45n term
- We say that the n<sup>2</sup> term is *dominant* in this expression

## **Big-Oh Notation**

 We wish a characterization of the time complexity of an algorithm that is *independent* on any implementation details (programming language and computer that will execute the algorithm).

# **Big-Oh Notation**

- The key issue is the *asymptotic complexity* of the function or *how it grows as n increases* 
  - This is determined by the *dominant term* in the growth function (the term that increases most quickly as *n* increases)
  - Constants become irrelevant as *n* increases since we want a characterization of the time complexity of an algorithm that is *independent* of the computer that will be used to execute it. Since different computers differ in speed by a constant factor, constant factors are ignored when expressing the asymptotic complexity of a function.

# **Big-Oh Notation**

- The asymptotic complexity of the function is referred to as the order of the function, and is specified by using *Big-Oh notation*.
  - Example: O(n<sup>2</sup>) means that the time taken by the algorithm grows like the n<sup>2</sup> function as n increases
  - O(1) means constant time, regardless of the size of the problem

### Some Growth Functions and Their Asymptotic Complexities

<b>Growth Function</b>	Order
t(n) = 17	O(1)
t(n) = 20n - 4	O(n)
t(n) = 12n * log <sub>2</sub> n + 100n	O(n*log <sub>2</sub> n)
t(n) = 3n <sup>2</sup> + 5n - 2	<b>O(n²)</b>
$t(n) = 2^n + 18n^2 + 3n$	<b>O(2</b> <sup>n</sup> )

#### Comparison of Some Typical Growth Functions



### **Exercise: Asymptotic Complexities**

<b>Growth Function</b>	Order
$t(n) = 5n^2 + 3n$	?
$t(n) = n^3 + \log_2 n - 4$	?
t(n) = log <sub>2</sub> n * 10n + 5	?
$t(n) = 3n^2 + 3n^3 + 3$	?
$t(n) = 2^n + 18n^{100}$	?

## **Determining Time Complexity**

- Algorithms frequently contain sections of code that are executed over and over again, i.e. *loops*
- Analysing loop execution is basic to determining time complexity

## Analysing Loop Execution

- A loop executes a certain number of times (say n), so the time complexity of the loop is n times the time complexity of the body of the loop
- Example: what is the time complexity of the following loop, in Big-O notation?

```
x = 0;
for (int i=0; i<n; i++)
x = x + 1;
```

- Nested loops: the body of the outer loop includes the inner loop
- **Example**: what is the time complexity of the following loop, in Big-O notation? Read the next set of notes from the course's webpage to see how the time complexity of this algorithm and the algorithms in the following pages are computed.

for (int i=0; i<n; i++) {

```
x = x + 1;
for (int j=0; j<n; j++)
y = y - 1;
}
```

### More Loop Analysis Examples

### More Loop Analysis Examples

### Analysis of Stack Operations

- Stack operations are generally efficient, because they all work on only one end of the collection
- But which is more efficient: the array implementation or the linked list implementation?

### Analysis of Stack Operations

- n is the number of items on the stack
- push operation for ArrayStack:
  - O(1) if array is not full (why?)
  - What would it be if the array is full? (worst case)
- push operation for LinkedStack:
  - O(1) (why?)
- pop operation for each?
- peek operation for each?