

Quantifier Elimination (QE)

Let $F(x_1, \dots, x_n)$ be a quantifier-free formula in disjunctive normal form. Let $P := (Q_{k+1}x_{k+1} \cdots Q_n x_n) F(x_1, \dots, x_n)$ be a prenex formula, where each Q_i is either the existential quantifier (\exists) or the universal quantifier (\forall). The goal of **Quantifier Elimination (QE)** is to compute a quantifier free formula $S(x_1, \dots, x_k)$ such that $SF \Leftrightarrow PF$ holds for all x_1, \dots, x_k .

As an example, consider

$$(\exists x)(\forall y)(ax^2 + bx + c) - (ay^2 + by + c) \geq 0$$

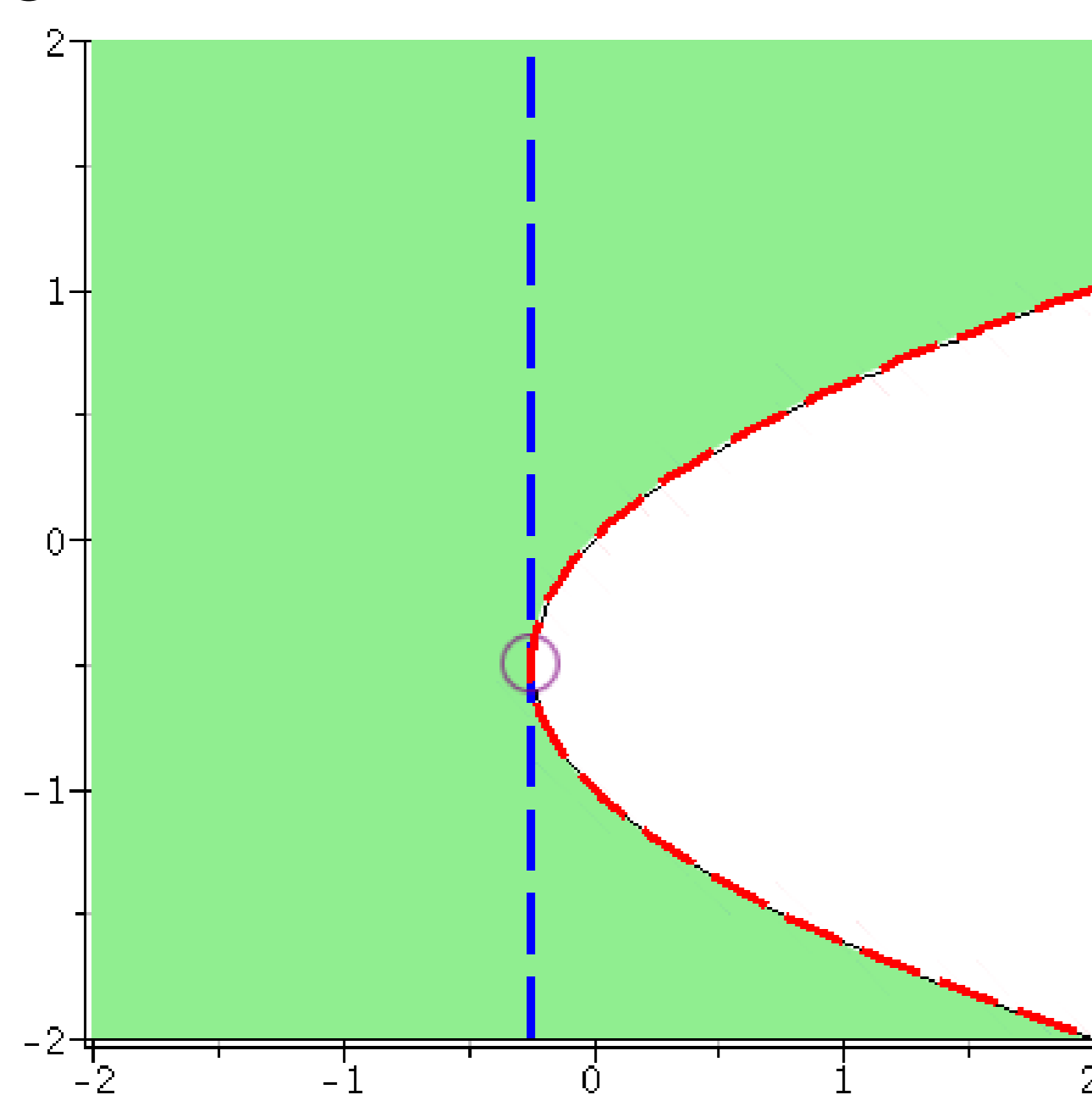
for which QE yields

$$(a < 0) \vee (a = b = 0).$$

In many research articles and computer software (QEPCAD, Mathematica, Reduce) QE is performed via cylindrical algebraic decomposition.

Cylindrical Algebraic Decomposition

A cylindrical algebraic decomposition (CAD) of \mathbb{R}^2 such that $y^2 + y - x$ is sign invariant in each of the 9 cells is shown below.



Motivation

In 1973, Collins introduced the first algorithm for computing CAD via a **projection and lifting** scheme. The projection phase relies on a **global projection operator**, which is a function defined **independently of the input system**, to eliminate variables. A strong projection operator usually produces much more polynomials than needed while a weak projection operator may fail for non-generic cases.

At ISSAC'09, the authors (with B. Xia and L. Yang) proposed a new scheme for computing CAD based on triangular decomposition (TCAD), which replaces the projection phase by computing a **complex cylindrical tree**. It is a natural question to ask how QE can be done via TCAD.

Main Contribution

We propose an operation which **refines a given complex cylindrical tree** with a new polynomial: this yields an **incremental** algorithm for computing the cylindrical decomposition of the complex space. An existing CAD can be refined through this operation. Moreover, QE is achieved based on **refined CADs**. This operation can **fully take advantage of equational constraints**.

Refining a Complex Cylindrical Tree

Fig.1 shows a $y^2 + x$ -invariant cylindrical decomposition of \mathbb{C}^2 , which is represented by a tree. This tree can be refined by the polynomial $y^2 + y$ into a $\{y^2 + x, y^2 + y\}$ -invariant cylindrical decomposition of \mathbb{C}^2 , see Fig. 2.

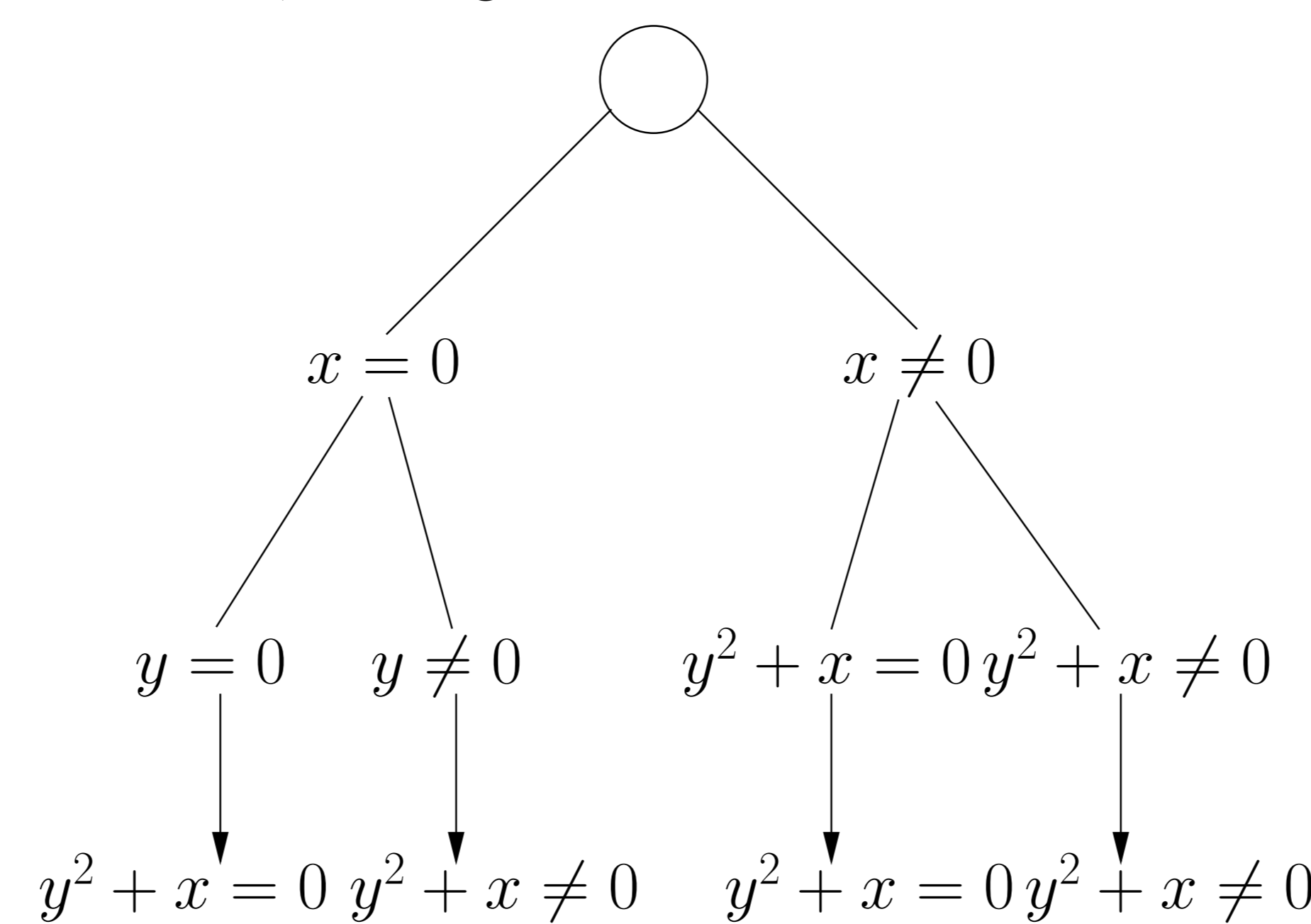


Fig.1

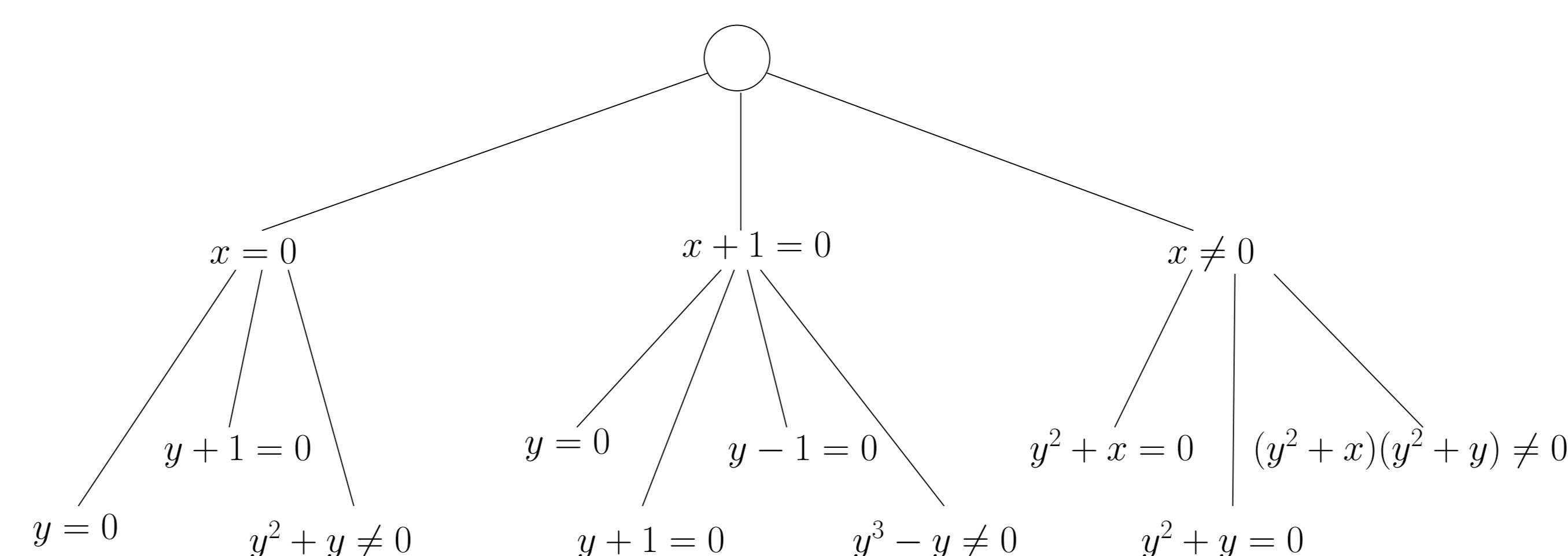


Fig. 2

Experimentation and Application

Results with a **preliminary implementation** in the MAPLE's **RegularChains** library.

Table 1. The timings for computing cylindrical decomposition of complex space

System	Recursive(t)	Incremental(t)	System	Recursive(t)	Incremental(t)
Amborg-Lazard-rev	2933.851	472.185	Lafferriere37	171.530	20.877
Barry	42.854	1.972	lhlp5	38.362	6.956
blood-coagulation-2	246.963	31.093	McCallum-random	> 1 hour	51.363
cdc2-cyclin	> 1 hour	120.923	nql-5-4	> 1 hour	9.620
circles	271.452	13.684	r-5	90.453	46.270
Collins-Johnson	20.405	4.464	r-6	1608.392	1151.847
Collision	899.116	59.347	Raksanyi	2177.672	961.672
Davenport-Heintz	12.172	2.076	Rose	> 1 hour	38.842
GonzalezGonzalez	166.182	105.590	Solotareff-3	315.375	23.525
hereman-2	> 1 hour	1574.710	X-axis-ellipse	145.505	35.358
Jirstrand22	62.095	10.448	xy-5-5-1	114.627	12.336
Jirstrand42	152.145	112.735	YangBaxterRosso	73.500	2.068

Table 2. The timings for computing CAD of real space and number of cells

System	TCAD(t)	TCAD(#)	QEPCAD(t)	QEPCAD(#)
Alonso	64.220	9939	12.060	39801
Barry	7.320	1087	Fail	Fail
blood-coagulation-2	1438.769	96093	Fail	Fail
cdc2-cyclin	801.470	435	> 1 hour	> 1 hour
Collins-Johnson	13.116	3673	10.272	3673
Collision	294.462	45979	11.880	45979
Davenport-Heintz	13.044	4949	9.868	4949
hereman-2	> 1 hour	> 1 hour	37.206	6899
Jirstrand41	2.716	755	9.76	755
Jirstrand42	> 1 hour	> 1 hour	26.761	179905
Lafferriere35	4.056	801	9.980	1005
Lafferriere37	102.450	13371	Fail	Fail
lhlp2	2.444	353	Fail	Fail
McCallum-random	> 1 hour	> 1 hour	28.585	877
nql-5-4	110.750	23347	94.273	23347
r-5	1312.090	261683	> 1 hour	> 1 hour
Rose	617.106	7951	Fail	Fail
Solotareff-3	> 1 hour	> 1 hour	14.028	66675

Consider a dynamical system with a control parameter u . We'd like to know if there exists a control Lyapunov function of the form $ax^2 + by^2$. The equivalent QE problem is: $\forall (x, y) \exists (u) (x \neq 0 \vee y \neq 0 \Rightarrow V > 0 \wedge \frac{d}{dt} V < 0)$.

The input control system.

$$\begin{aligned} > S := \{\text{diff}(x(t), t) = -x(t)+u, \text{diff}(y(t), t) = -x(t) - y(t)^3\}; \\ S := \left\{ \frac{d}{dt} x(t) = -x(t) + u, \frac{d}{dt} y(t) = -x(t) - y(t)^3 \right\} \end{aligned}$$

Compute $\frac{d}{dt} V$.

$$\begin{aligned} > fx := -x+u; fy := -x-y^3; V := a*x^2+b*y^2; \\ Vt := \text{diff}(V, x)*fx + \text{diff}(V, y)*fy; \\ \begin{aligned} fx &:= -x + u \\ fy &:= -x - y^3 \\ V &:= y^2 b + x^2 a \\ Vt &:= 2 x a (-x + u) + 2 b y (-x - y^3) \end{aligned} \end{aligned}$$

Call QeTcad to find conditions on a and b .

$$> \text{RegularChains:-QuantifierElimination}([\{A, x, y\}, [E, u]], [\{x \neq 0\}, \{y \neq 0\}], [\{V > 0, Vt < 0\}]);$$

Verify the result for $a = 1, b = 4$. Indeed $x^2 + 4y^2$ is a control Lyapunov function.

$$\begin{aligned} > \text{simplify}(\text{subs}(\{a=1, b=4\}, Vt)); \text{subs}(u=4*y, \%); \\ \begin{aligned} 2 u x - 2 x^2 - 8 y^4 - 8 y x \\ - 2 x^2 - 8 y^4 \end{aligned} \end{aligned}$$