

Exact Computation of the Real Solutions of Arbitrary Polynomial Systems

Presented by Marc Moreno Maza¹

joint work with

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Plan

- 1 Computing the real solutions of polynomial systems symbolically
- 2 Triangular decomposition of semi-algebraic systems
- 3 Application to dynamical system analysis

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Solving polynomial systems? What does this mean?

The algebra text book says:

For $F \subset \mathbf{k}[x_1, \dots, x_n]$ this is simply

- a *primary decomposition* of $\langle F \rangle$ or
- the *irreducible decomposition* of $V(F)$ (the zero set of F in $\bar{\mathbf{k}}^n$).

The computer algebra system does well:

For $F \subset \mathbf{k}[x_1, \dots, x_n]$, with $\mathbf{k} = \mathbb{Z}/p\mathbb{Z}$ or $\mathbf{k} = \mathbb{Q}$,

- computing a *Gröbner basis* of $\langle F \rangle$ or
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But most scientists and engineers need:

- For $F \subset \mathbb{Q}[x_1, \dots, x_n]$, a useful description of the points of $V(F)$ whose coordinates are real.
- For $F \subset \mathbb{Q}[u_1, \dots, u_d][x_1, \dots, x_n]$, the real (x_1, \dots, x_n) -solutions as a function of the real parameter (u_1, \dots, u_d) .

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Solving for the real solutions: classical techniques

In dimension zero over \mathbb{Q} :

For $F \subset \mathbb{Q}[x_1, \dots, x_n]$, if $V(F)$ is finite, many standard and efficient techniques apply to identify the real solutions.

In (generic) dimension zero over $\mathbb{Q}[u_1, \dots, u_d]$:

For $F \subset \mathbb{Q}[u_1, \dots, u_d][x_1, \dots, x_n]$ and an integer r one can determine “generic” conditions on u_1, \dots, u_d for F to admit exactly r real (x_1, \dots, x_n) -solutions

For arbitrary systems:

For $F \subset \mathbb{Q}[x_1, \dots, x_n]$, one can partition \mathbb{R}^n into *cylindrical cells* where the sign of each $f \in F$ does not change.

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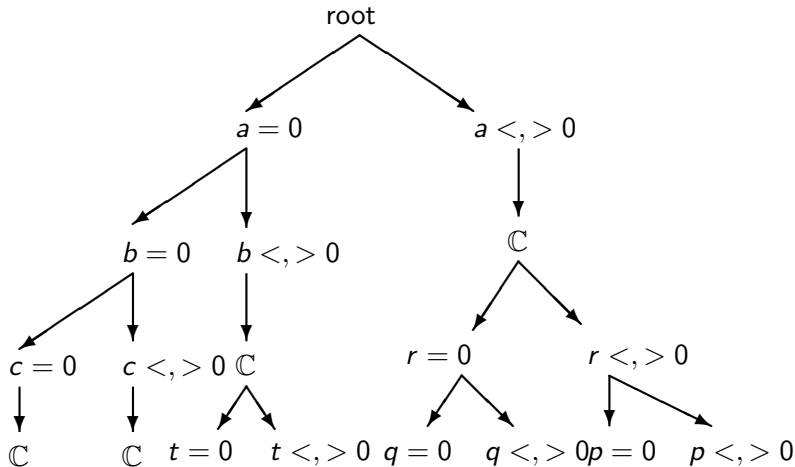
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Cylindrical algebraic decomposition of $\{ax^2 + bx + c\}$



The cylindrical algebraic decomposition of $\{ax^2 + bx + c\}$ is given by the tree above, where $t = bx + c$, $q = 2ax + b$, and $r = 4ac - b^2$. This is the best possible output for that method, leading to **27 cells!**

Can a computer program be as good as a high-school student?

For the equation $ax^2 + bx + c = 0$, can a computer program produce?

$$\left\{ \begin{array}{l} ax^2 + bx + c = 0 \\ a \neq 0 \wedge b^2 - 4ac > 0 \end{array} \right. \quad \left\{ \begin{array}{l} 2ax + b = 0 \\ 4ac - b^2 = 0 \\ a \neq 0 \end{array} \right.$$

$$\left\{ \begin{array}{l} bx + c = 0 \\ a = 0 \\ b \neq 0 \end{array} \right. \quad \left\{ \begin{array}{l} c = 0 \\ b = 0 \\ a = 0 \end{array} \right.$$

Yes, our new algorithm RealTriangularize can do that!

```
Applications Places System [Icons] Untitled (1) - [Server 1] - Maple 14 moreno [Icons]
```

```
with(RegularChains) : with(SemiAlgebraicSetTools) : with(ParametricSystemTools) : with(ParametricSystemTools)
```

```
R := PolynomialRing([x, c, b, a]); F := [a·x2 + b·x + c];  
polynomial_ring  
[a x2 + b x + c]
```

Solving for the real solutions:

```
RealTriangularize(F, R, output = record);
```

$$\left[\begin{array}{l} a x^2 + b x + c = 0 \\ -4 c a + b^2 > 0 \\ a \neq 0 \end{array} \right], \left[\begin{array}{l} b x + c = 0 \\ b \neq 0 \\ a = 0 \end{array} \right], \left[\begin{array}{l} c = 0 \\ b = 0 \\ a = 0 \end{array} \right], \left[\begin{array}{l} 2 a x + b = 0 \\ 4 a c - b^2 = 0 \\ a \neq 0 \end{array} \right]$$

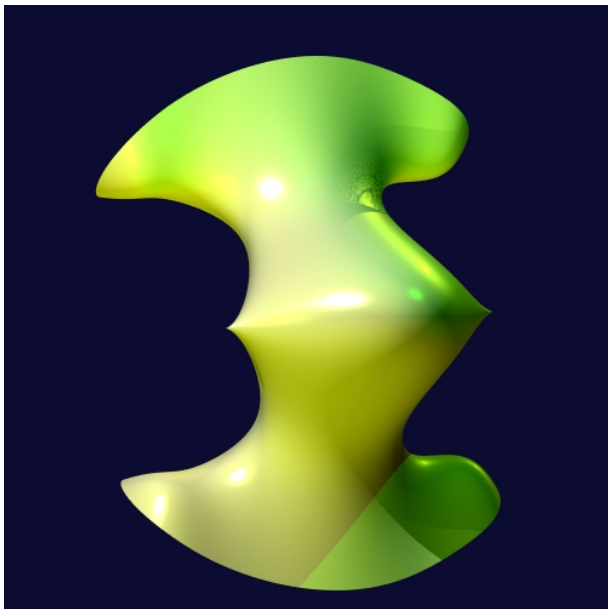
Solving for the complex solutions

```
dec := Triangularize(F, R, output = lazard); map(Display, dec, R);
```

```
[regular_chain, regular_chain, regular_chain]
```

$$\left[\left[\begin{array}{l} a x^2 + b x + c = 0 \\ a \neq 0 \end{array} \right], \left[\begin{array}{l} b x + c = 0 \\ a = 0 \\ b \neq 0 \end{array} \right], \left[\begin{array}{l} c = 0 \\ b = 0 \\ a = 0 \end{array} \right] \right]$$

RealTriangularize applied to the *Eve* surface (1/2)



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Regular chain

Example

$$T := \begin{cases} f_3 = x_4 x_5^2 + 2x_5 + 1 \\ f_2 = (x_1 + x_2) x_3^2 + x_3 + 1 \\ f_1 = x_1^2 - 1. \end{cases}$$

Under the order $x_5 > x_4 > x_3 > x_2 > x_1$,

- $\text{mvar}(f_2) = x_3$ and $\text{mvar}(f_1) = x_1$
- $\text{init}(f_2) = x_1 + x_2$ and $\text{init}(f_1) = 1$
- T is a regular chain. Indeed:
 - ▶ $\text{init}(f_2)$ is regular (neither zero nor zero-divisor) modulo $\langle f_1 \rangle$.
 - ▶ $\text{init}(f_3)$ is regular modulo $\langle f_1, f_2 \rangle : \text{init}(f_2)^\infty$.

Proposition

Let $T = \{t_1, \dots, t_s\}$ be a regular chain of $\mathbf{k}[x]$. Then the saturated ideal $\text{sat}(T) := \langle T \rangle : (\prod_{i=1}^s \text{init}(t_i))^\infty$ is a proper equi-dimensional ideal of $\mathbf{k}[x]$.

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Regular semi-algebraic system

Notation

- Let $T \subset \mathbb{Q}[x_1 < \dots < x_n]$ be a regular chain with $\mathbf{y} := \{\text{mvar}(t) \mid t \in T\}$ and $\mathbf{u} := \mathbf{x} \setminus \mathbf{y} = u_1, \dots, u_d$.
- Let P be a finite set of polynomials, s.t. every $f \in P$ is regular modulo $\text{sat}(T)$.
- Let Q be a quantifier-free formula of $\mathbb{Q}[\mathbf{u}]$.

Definition

We say that $R := [Q, T, P_{>}]$ is a **regular semi-algebraic system** if:

- Q defines a **non-empty open** semi-algebraic set S in \mathbb{R}^d ,
- the regular system $[T, P]$ **specializes well** at every point u of S
- at each point u of S , the specialized system $[T(u), P(u)_{>}]$ has **at least one real solution**.

$$Z_{\mathbb{R}}(R) = \{(u, y) \mid Q(u), t(u, y) = 0, p(u, y) > 0, \forall (t, p) \in T \times P\}.$$

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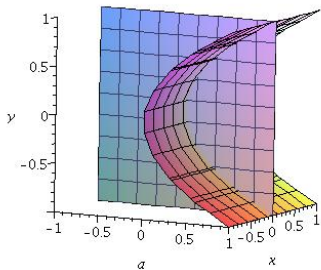
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Example

The system $[Q, T, P_{>}]$, where

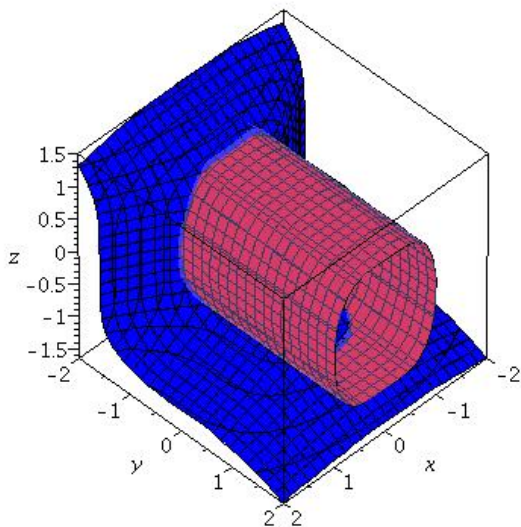
$$Q := a > 0, \quad T := \begin{cases} y^2 - a = 0 \\ x = 0 \end{cases}, \quad P_{>} := \{y > 0\}$$

is a regular semi-algebraic system.



RealTriangularize applied to *sofa* and *cylinder* (1/2)

$$x^2 + y^3 + z^5 = x^4 + z^2 - 1 = 0$$



RealTriangularize applied to *sofa* and *cylinder* (2/2)

Applications Places System

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true

```
> R := PolynomialRing([z, y, x]); F := [x^2 + y^3 + z^5, x^4 + z^2 - 1]; RealTriangularize(F, R, output =
      R := polynomial_ring
      F := [x^2 + y^3 + z^5, x^4 + z^2 - 1]
```

$$\left\{ \begin{array}{l} (-2x^4 + x^8 + 1)z + y^3 + x^2 = 0 \\ y^6 + 2x^2y^3 + 10x^{12} - 10x^8 + x^{20} - 5x^{16} + 6x^4 - 1 = 0 \\ x < 1 \\ x + 1 > 0 \\ x^{12} - 4x^8 + 5x^4 - 1 \neq 0 \end{array} \right. , \left\{ \begin{array}{l} (-2x^4 + x^8 + 1)z - x^2 = 0 \\ y^3 + 2x^2 = 0 \\ x^{12} - 4x^8 + 5x^4 - 1 = 0 \end{array} \right. ,$$

$$\left\{ \begin{array}{l} (-2x^4 + x^8 + 1)z + x^2 = 0 \\ y = 0 \\ x^{12} - 4x^8 + 5x^4 - 1 = 0 \end{array} \right. , \left\{ \begin{array}{l} z = 0 \\ y + 1 = 0 \\ x + 1 = 0 \end{array} \right. , \left\{ \begin{array}{l} z = 0 \\ y + 1 = 0 \\ x - 1 = 0 \end{array} \right.$$

```
> |
```


The RegularChains library in MAPLE (1/2)

Design goals

- Solving polynomial systems over \mathbb{Q} and \mathbb{F}_p , including **parametric** systems and **semi-algebraic** systems.
- Offering tools to manipulate their solutions.
- Organized around the concept of a **regular chain**, accommodating all types of solving and providing space-and-time efficiency.

Features

- Use of types for algebraic structures: `polynomial_ring`, `regular_chain`, `constructible_set`, `quantifier_free_formula`, `regular_semi_algebraic_system`.
- Top level commands: `PolynomialRing`, `Triangularize`, `RealTriangularize` `SamplePoints`, ...
- Tool kits: `ConstructibleSetTools`, `ParametricSystemTools`, `FastArithmeticTools`, `SemiAlgebraicSetTools`, ...

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The RegularChains library in MAPLE (2/2)

Classical tools

- Isolating the real solutions of zero-dimensional polynomial systems:
`SemiAlgebraicSetTools:-RealRootIsolate`
- Real root classification of parametric polynomial systems:
`ParametricSystemTools:-RealRootClassification`
- Cylindrical algebraic decomposition of polynomial systems:
`SemiAlgebraicSetTools:-CylindricalAlgebraicDecompose`

New tools

- Triangular decomposition of semi-algebraic systems:
`RealTriangularize`
- Sampling all connected components of a semi-algebraic system:
`SamplePoints`
- Set-theoretical operations on semi-algebraic sets:
`SemiAlgebraicSetTools:-Difference`

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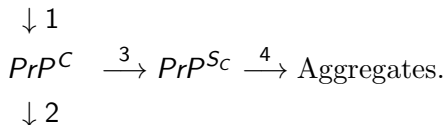
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A biochemical network: setting (1/3)

The generic kinetic scheme of prion diseases is illustrated as follows:

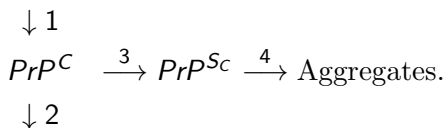


where

- $[PrP^C]$ is the concentration of PrP^C (harmless form)
- $[PrP^{Sc}]$ is the concentration of PrP^{Sc} (infectious form) which catalyses the transformation from the normal form to itself,
- **Step 1:** synthesis of native PrP^C
- **Step 2, 4:** normal degradation.

A biochemical network: setting (2/3)

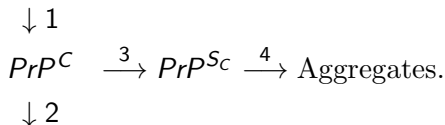
Let ν_i be the rate of Step i for $i = 1, \dots, 4$.



- **Step 1:** zero-order kinetic process, that is $\nu_1 = k_1$,
- **Step 2, 4:** first-order rate equations: $\nu_2 = k_2 [PrP^C]$,
 $\nu_4 = k_4 [PrP^{Sc}]$.
- **Step 3:** a nonlinear process

$$\nu_3 = [PrP^C] \frac{a(1 + b[PrP^{Sc}]^n)}{1 + c[PrP^{Sc}]^n}.$$

A biochemical network: setting (3/3)



We also have:

$$\begin{aligned} \frac{d [PrP^C]}{dt} &= \nu_1 - \nu_2 - \nu_3 \\ \frac{d [PrP^{Sc}]}{dt} &= \nu_3 - \nu_4 \end{aligned}$$

Letting $x = [PrP^C]$ and $y = [PrP^{Sc}]$, we obtain the dynamical system:

$$\begin{aligned} \frac{dx}{dt} &= k_1 - k_2x - ax \frac{(1 + by^n)}{1 + cy^n} \\ \frac{dy}{dt} &= ax \frac{(1 + by^n)}{1 + cy^n} - k_4y \end{aligned}$$

A biochemical network: semi-algebraic systems to solve

Dynamical system to study

M. Laurent (Biochem. J., 1996) suggests to set $b = 2$, $c = 1/20$, $n = 4$, $a = 1/10$, $k_4 = 50$ and $k_1 = 800$, leading to:

$$\begin{cases} \frac{dx}{dt} = f_1 \\ \frac{dy}{dt} = f_2 \end{cases} \quad \text{with} \quad \begin{cases} f_1 = \frac{16000 + 800y^4 - 20k_2x - k_2xy^4 - 2x - 4xy^4}{20 + y^4} \\ f_2 = \frac{2(x + 2xy^4 - 500y - 25y^5)}{20 + y^4} \end{cases} .$$

Semi-algebraic systems to solve

By Routh-Hurwitz criterion, an equilibrium (x, y) is asymptotically stable if

$$\Delta_1 := -\left(\frac{\partial f_1}{\partial x} + \frac{\partial f_2}{\partial y}\right) > 0 \quad \text{and} \quad a_2 := \frac{\partial f_1}{\partial x} \cdot \frac{\partial f_2}{\partial y} - \frac{\partial f_1}{\partial y} \cdot \frac{\partial f_2}{\partial x} > 0.$$

Letting p_1, p_2 the above polynomials, we obtain two semi-algebraic systems:

$$\mathcal{S}_1 : \{p_1 = p_2 = 0, k_2 > 0\} \quad \text{and} \quad \mathcal{S}_2 : \{p_1 = p_2 = 0, k_2 > 0, \Delta_1 > 0, a_2 > 0\}$$

A biochemical network: solving \mathcal{S}_1

The real solutions of \mathcal{S}_1 are described by the following **triangular decomposition** into **regular semi-algebraic systems**.

$$A_1 := \begin{cases} (2y^4 + 1)x - 25y^5 - 500y & = 0 \\ (k_2 + 4)y^5 - 64y^4 + (2 + 20k_2)y - 32 & = 0 \\ k_2 & > 0 \\ r & \neq 0 \end{cases}, \quad A_2 := \begin{cases} t_x & = 0 \\ t_y & = 0 \\ r & = 0 \\ k_2 & > 0 \end{cases}.$$

where $t_y(k_2, y)$ has degree 4 in y and r is given by

$$r := 100000k_2^8 + 1250000k_2^7 + 5410000k_2^6 + 8921000k_2^5 - 9161219950k_2^4 \\ - 5038824999k_2^3 - 1665203348k_2^2 - 882897744k_2 + 1099528405056.$$

The polynomial r has four real roots, two are positive: $\alpha_1 < \alpha_2$.

A biochemical network: solving \mathcal{S}_1

Regarding k_2 as a parameter, one can compute a **real comprehensive triangular decomposition** which gives:

$$\left\{ \begin{array}{ll} \{ \} & k_2 \leq 0 \\ \{A_1\} & 0 < k_2 < \alpha_1 \\ \{A_2\} & k_2 = \alpha_1 \\ \{A_1\} & \alpha_1 < k_2 < \alpha_2 \\ \{A_2\} & k_2 = \alpha_2 \\ \{A_1\} & k_2 > \alpha_2 \end{array} \right. .$$

From where we deduce the number of real solutions:

$$\left\{ \begin{array}{ll} 0 & k_2 \leq 0 \\ 1 & k_2 > 0 \text{ and } r > 0 \\ 2 & k_2 > 0 \text{ and } r = 0 \\ 3 & k_2 > 0 \text{ and } r < 0 \end{array} \right.$$

A biochemical network: conclusion

Theorem

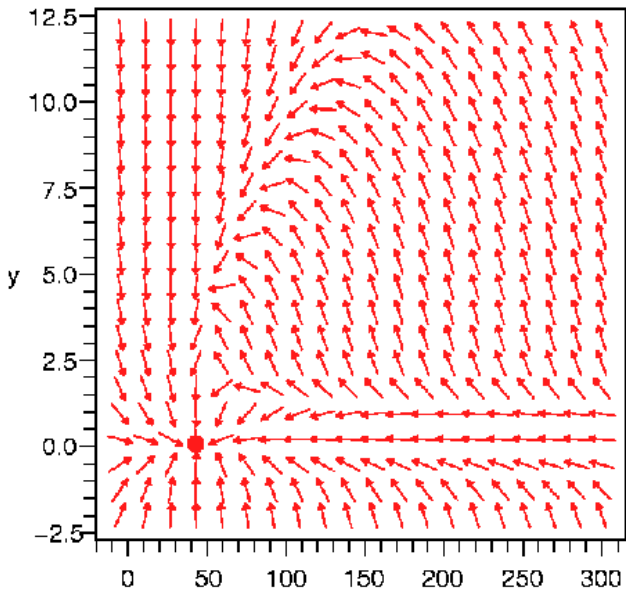
Assume that $k_2 > 0$. Then we have: if $r > 0$, then the dynamical system has 1 equilibrium; if $r = 0$, then it has 2 equilibria; if $r < 0$, it has 3 equilibria.

Theorem

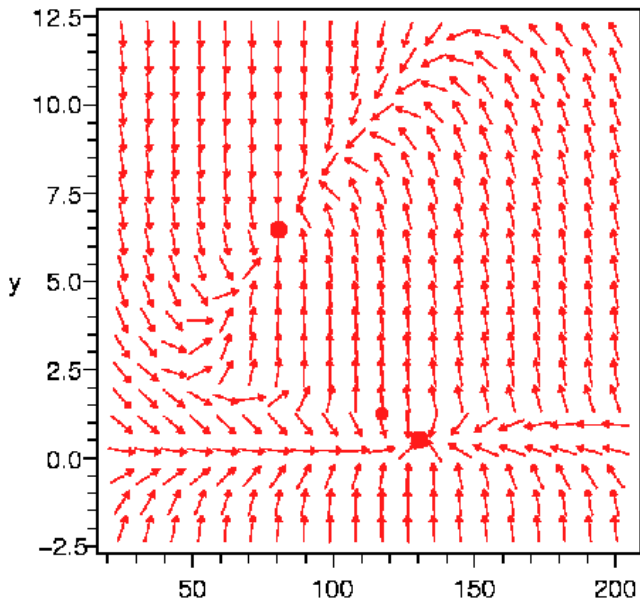
Assume that $k_2 > 0$. Then we have: if $r > 0$, then the system has one hyperbolic equilibrium, which is asymptotically stable; if $r < 0$ and $r_2 \neq 0$, then the system has three hyperbolic equilibria, two of which are asymptotically stable and the other one is unstable; if $r = 0$ or $r_2 = 0$, the system experiences a bifurcation where

$$\begin{aligned} r_2 = & 10004737927168k_2^9 + 624166300700672k_2^8 + 7000539052537600k_2^7 \\ & + 45135589467012800k_2^6 - 840351411856453750k_2^5 - 50098004352248446875k_2^4 \\ & - 27388168989455000000k_2^3 - 8675209266696000000k_2^2 \\ & + 102960917356800000000k_2 + 5932546064102400000000. \end{aligned}$$

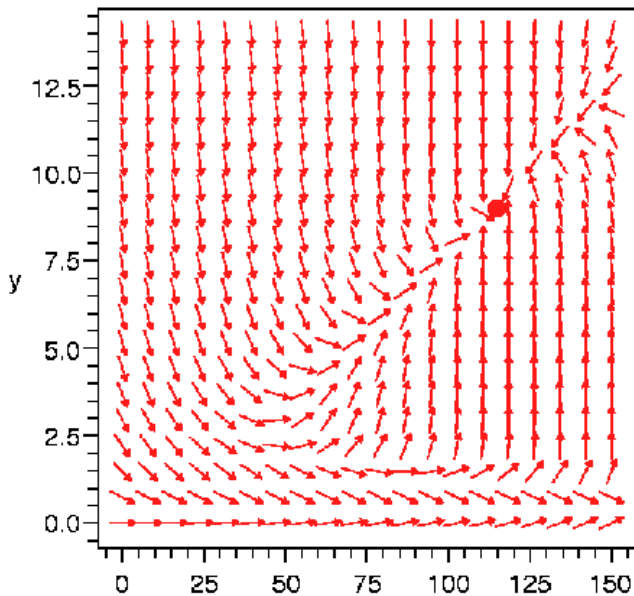
k2=18



$k_2 = 8$



$k^2 = 3$



Summary and notes

- Solving for the **real roots** of (parametric or not) polynomial systems is a fundamental problem with many applications.
- Most of the time, this requires **exact (thus symbolic) computation**.
- Computer algebra systems used to have limited capabilities for that, especially for the parametric case.
- Recent work (Changbo Chen, James H. Davenport, M^3 , Bican Xia & Rong Xiao, ISSAC 2010-2011) is changing that.
- **RealTriangularize** is available MAPLE 15, as part of the MAPLE **RegularChains** library.