

# Computing the real solutions of polynomial systems with the RegularChains library in MAPLE

Presented by Marc Moreno Maza<sup>1</sup>

joint work with

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# Plan

- 1 Overview
- 2 Solver verification
- 3 Branch cut analysis
- 4 Biochemical network analysis
- 5 Reachability problem for hybrid systems

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# The RegularChains library in MAPLE

## Design goals

- Solving polynomial systems over  $\mathbb{Q}$  and  $\mathbb{F}_p$ , including **parametric** systems and **semi-algebraic** systems.
- Offering tools to manipulate their solutions.
- Organized around the concept of a **regular chain**, accommodating all types of solving and providing space-and-time efficiency.

## Features

- Use of types for algebraic structures: `polynomial_ring`, `regular_chain`, `constructible_set`, `quantifier_free_formula`, `regular_semi_algebraic_system`.
- Top level commands: `PolynomialRing`, `Triangularize`, `RealTriangularize` `SamplePoints`, ...
- Tool kits: `ConstructibleSetTools`, `ParametricSystemTools`, `FastArithmeticTools`, `SemiAlgebraicSetTools`, ...

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# Solving for the real solutions of polynomial systems

## Classical tools

- Isolating the real solutions of zero-dimensional polynomial systems:  
`SemiAlgebraicSetTools:-RealRootIsolate`
- Real root classification of parametric polynomial systems:  
`ParametricSystemTools:-RealRootClassification`
- Cylindrical algebraic decomposition of polynomial systems:  
`SemiAlgebraicSetTools:-CylindricalAlgebraicDecompose`

## New tools

- Triangular decomposition of semi-algebraic systems:  
`RealTriangularize`
- Sampling all connected components of a semi-algebraic system:  
`SamplePoints`
- Set-theoretical operations on semi-algebraic sets:  
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# Regular semi-algebraic system

## Notation

- Let  $T \subset \mathbb{Q}[x_1 < \dots < x_n]$  be a regular chain with  $\mathbf{y} := \{\text{mvar}(t) \mid t \in T\}$  and  $\mathbf{u} := \mathbf{x} \setminus \mathbf{y} = u_1, \dots, u_d$ .
- Let  $P$  be a finite set of polynomials, s.t. every  $f \in P$  is regular modulo  $\text{sat}(T)$ .
- Let  $Q$  be a quantifier-free formula of  $\mathbb{Q}[\mathbf{u}]$ .

## Definition

We say that  $R := [Q, T, P_{>}]$  is a **regular semi-algebraic system** if:

- $Q$  defines a **non-empty open** semi-algebraic set  $S$  in  $\mathbb{R}^d$ ,
- the regular system  $[T, P]$  **specializes well** at every point  $u$  of  $S$
- at each point  $u$  of  $S$ , the specialized system  $[T(u), P(u)_{>}]$  has **at least one real solution**.

$$Z_{\mathbb{R}}(R) = \{(u, y) \mid Q(u), t(u, y) = 0, p(u, y) > 0, \forall (t, p) \in T \times P\}.$$



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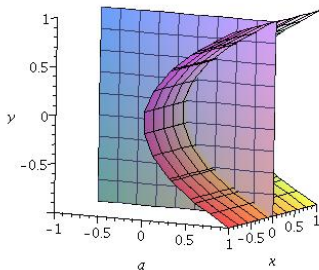
$$Z_{\mathbb{R}}(R) = \{(u, y) \mid Q(u), t(u, y) = 0, p(u, y) > 0, \forall (t, p) \in T \times P\}.$$

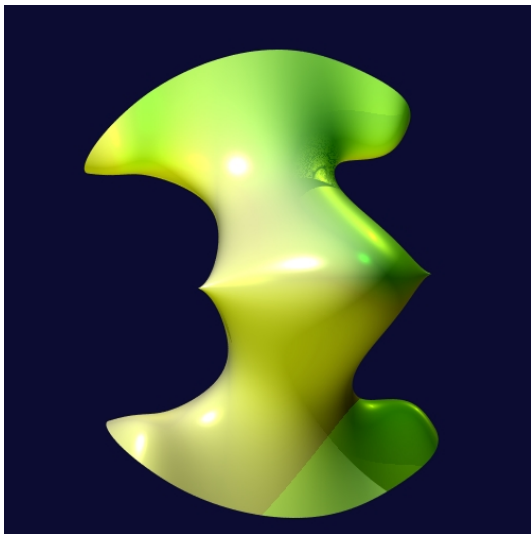
## Example

The system  $[Q, T, P_{>}]$ , where

$$Q := a > 0, \quad T := \begin{cases} y^2 - a = 0 \\ x = 0 \end{cases}, \quad P_{>} := \{y > 0\}$$

is a regular semi-algebraic system.



RealTriangularize applied to the *Eve* surface (1/2)

RealTriangularize applied to the *Eve* surface (2/2)

```

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Format Table Drawing Plot Spreadsheet Tools Window Help

```

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R := PolynomialRing([x, y, z]); F := [5*x^2 + 2*x*z^2 + 5*y^6 + 15*y^4 + 5*z^2 - 15*y^5 - 5*y^3];
                                     polynomial_ring

```

$$[5x^2 + 2xz^2 + 5y^6 + 15y^4 + 5z^2 - 15y^5 - 5y^3]$$

```

RealTriangularize(F, R, output = record);

```

$$\begin{cases} 5x^2 + 2z^2x + 5y^6 + 15y^4 - 5y^3 - 15y^5 + 5z^2 = 0 \\ 25y^6 - 75y^5 + 75y^4 - z^4 - 25y^3 + 25z^2 < 0 \end{cases}$$

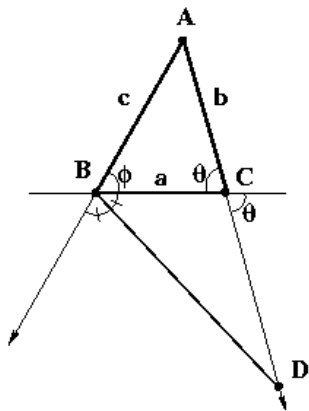
$$\begin{cases} 5x + z^2 = 0 \\ 25y^6 - 75y^5 + 75y^4 - 25y^3 - z^4 + 25z^2 = 0 \\ 64z^4 - 1600z^2 + 25 > 0 \\ z \neq 0 \\ z - 5 \neq 0 \\ z + 5 \neq 0 \end{cases}, \begin{cases} x = 0 \\ y - 1 = 0 \\ z = 0 \end{cases}, \begin{cases} x = 0 \\ y = 0 \\ z = 0 \end{cases}, \begin{cases} x + 5 = 0 \\ y - 1 = 0 \\ z - 5 = 0 \end{cases}$$

$$\begin{cases} x + 5 = 0 \\ y = 0 \\ z - 5 = 0 \end{cases}, \begin{cases} x + 5 = 0 \\ y - 1 = 0 \\ z + 5 = 0 \end{cases}, \begin{cases} x + 5 = 0 \\ y = 0 \\ z + 5 = 0 \end{cases}, \begin{cases} 5x + z^2 = 0 \\ 2y - 1 = 0 \\ 64z^4 - 1600z^2 + 25 = 0 \end{cases}$$

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# Are these two different output equivalent?



Given a triangle with edge lengths  $a, b, c$  (denoting the respective edges  $a, b, c$  too) the following two conditions  $S_1, S_2$  are both characterizing the fact that the external bisector of the angle of  $a, c$  intersects with  $b$  on the other side of  $a$  than the triangle:

$$S_1 = a > 0 \wedge b > 0 \wedge c > 0 \wedge a < b + c \wedge b < a + c \wedge c < a + b \wedge (b^2 + a^2 - c^2 \leq 0) \vee (c(b^2 + a^2 - c^2)^2 < ab^2(2ac - (c^2 + a^2 - b^2))),$$

$$S_2 = a > 0 \wedge b > 0 \wedge c > 0 \wedge a < b + c \wedge b < a + c \wedge c < a + b \wedge c - a > 0.$$

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## Is this simplification correct?

### The original problem

The branch cut of  $\sqrt{z}$  is conventionally:

$$\{z \in \mathbb{C} \mid \Re(z) < 0 \wedge \Im(z) = 0\}.$$

Do the following equations hold for all  $z \in \mathbb{C}$ :

$$\sqrt{z-1}\sqrt{z+1} = \sqrt{z^2-1} \quad \text{and} \quad \sqrt{1-z}\sqrt{1+z} = \sqrt{1-z^2}.$$

### Turning the question to sampling

- The branch cuts of each formula is a semi-algebraic system  $\mathcal{S}$  given as the disjunction of 3 others  $\mathcal{S}_1, \mathcal{S}_2, \mathcal{S}_3$  (one per  $\sqrt{\quad}$ ).
- Consider CAD-cells  $C_1, \dots, C_e$ , forming an intersection-free basis refining the connected components of  $\mathcal{S}_1, \mathcal{S}_2, \mathcal{S}_3$ .
- By virtue of the *Modromy Theorem*, it is sufficient to check whether the formula holds at a sample point of each of  $C_1, \dots, C_e$ .



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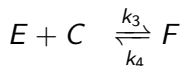
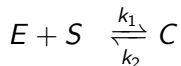
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# Is there a unique positive equilibrium?

*Allosteric enzyme*



$$\frac{1}{2} C - \frac{1}{2} E + S - \frac{1}{2} C_0 + \frac{1}{2} E_0 - S_0 = 0$$

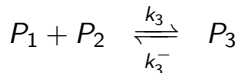
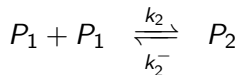
$$\frac{1}{2} C + \frac{1}{2} E + F - \frac{1}{2} C_0 - \frac{1}{2} E_0 - F_0 = 0$$

$$k_1 ES - k_2 C - k_3 EC + k_4 F = 0$$

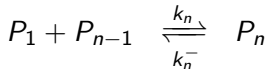
$$-2k_3 EC + 2k_4 F = 0$$

$$E, S, C, F, E_0, S_0, C_0, F_0, k_1, k_2, k_3, k_4 > 0.$$

*Cascad of polymerisation*



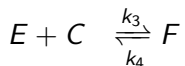
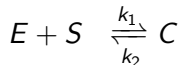
$\vdots$



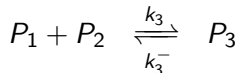
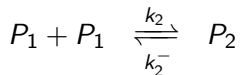
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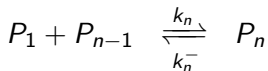
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$$\frac{1}{2} C + \frac{1}{2} E + F - \frac{1}{2} C_0 - \frac{1}{2} E_0 - F_0 = 0$$

$$k_1 ES - k_2 C - k_3 EC + k_4 F = 0$$

$$-2 k_3 EC + 2 k_4 F = 0$$

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## Which states can be reached from the initial one?

### Hybrid systems with linear control

- Given  $A \in \mathbb{R}^{n \times n}$ ,  $B \in \mathbb{R}^{n \times m}$  consider  $\dot{\xi} = A\xi + Bu$  where  $\xi(t) \in \mathbb{R}^n$  is the **state of the system** at time  $t$  and  $u : \mathbb{R} \rightarrow \mathbb{R}^m$  is a piecewise continuous function which is called the **control input**.
- With  $x = \xi(0)$  and a control input  $u$ , we have:
 
$$\xi(t) = \Phi(x, u, t) = e^{At}x + \int_0^t e^{A(t-\tau)}Bu(\tau)d\tau.$$
- Question: Given  $x = \xi(0)$ , which values  $\xi(t)$  can be reached?

### The (Lafferriere et al. 2001) example

$$\Phi(x_1, x_2, u, t) = \left( x_1 e^{2t} + \frac{2}{3}u(-e^{2t} + e^{\frac{1}{2}t}), x_2 e^{-t} + \frac{1}{2}u(e^t - e^{-t}) \right).$$

Let  $z = e^{\frac{1}{2}t}$ , the problem reduces to compute the  $(y_1, y_2)$  such that:

$$\exists u \exists z (0 \leq u \wedge z \geq 1 \wedge p_1 = 0 \wedge p_2 = 0) \text{ where}$$

$$p_1 = y_1 - \frac{2}{3}u(-z^4 + z) \text{ and } p_2 = y_2 z^2 - \frac{1}{2}u(z^4 - 1).$$

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