

Computing Cylindrical Algebraic Decomposition via Triangular Decomposition

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Background

Cylindrical algebraic decomposition (CAD) is a fundamental tool in real algebraic geometry. It was introduced by Collins in 1973 and has been followed by lots of improvements, like

- ▶ improved projection methods
(McCallum 88, 98, Hong 90, Brown 01)
- ▶ partially built CADs
(Collins and Hong 91, McCallum 93, Strzeboński 00)
- ▶ improved stack construction
(Collins, Johnson and Krandick 02)
- ▶ efficient projection orders
(Dolzmann, Seidl and Sturm 04)
- ▶ ...

Motivation

1. Understand the **relations and possible interactions** between CAD and triangular decompositions of polynomial systems.

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1. Understand the **relations and possible interactions** between CAD and triangular decompositions of polynomial systems.
2. Investigate the possibility of improving the **practical efficiency** of CAD implementation by means of modular methods and fast polynomial arithmetic, being developed for triangular decompositions.

Cylindrical Algebraic Decomposition (I)

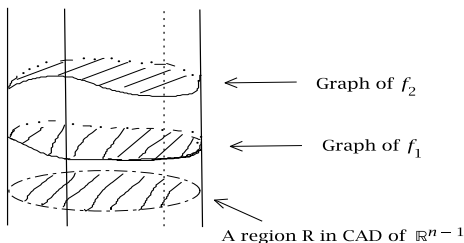
A cylindrical algebraic decomposition of \mathbb{R}^n can be defined inductively as follows.

- $n = 1$. A CAD of \mathbb{R} is a finite partition of the real line into points and open intervals.

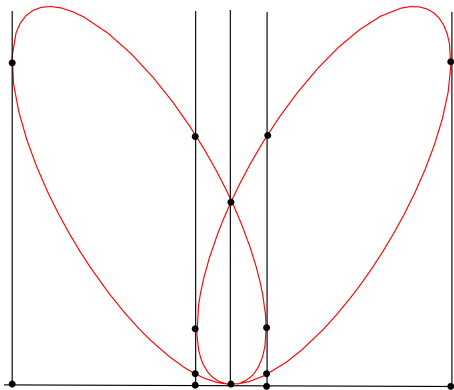


Cylindrical Algebraic Decomposition (II)

- $n > 1$. Given a CAD D' of \mathbb{R}^{n-1} , one builds a CAD D of \mathbb{R}^n as follows. Above each region R of D' :
 - ▶ consider finitely many disjoint graphs (called *sections*) of continuous real-valued algebraic functions,
 - ▶ decomposing the cylinder $R \times \mathbb{R}^1$, into sections and *sectors* (located between two consecutive sections), which form a *stack* over R ,
 - ▶ then all the sections and sectors are the elements of D .



A Cylindrical Algebraic Decomposition of \mathbb{R}^2 Induced by the Tacnode Curve



Tacnode curve: $y^4 - 2y^3 + y^2 - 3x^2y + 2x^4 = 0$.

Algorithm of Collins

Projection: Starting from the input $F_n \subset \mathbb{Q}[y_1, \dots, y_n]$, repeatedly apply a projection operator to eliminate the variables one by one until a set of univariate polynomials are obtained

$$F_n \rightarrow F_{n-1} \rightarrow \dots \rightarrow F_1$$

such that an F_k -invariant CAD of \mathbb{R}^k can be constructed from an F_{k-1} -invariant CAD of \mathbb{R}^{k-1} , for $2 \leq k \leq n$.

Lifting: One isolates the real roots of polynomial in F_1 and deduces a CAD of \mathbb{R}^1 . For each region of the CAD of \mathbb{R}^1 , one evaluates the polynomials of F_2 at a *sample point* and isolates their real roots, from which one produces a stack over the region. Continuing in this manner, one finally obtains a CAD of \mathbb{R}^n .

Another View of CAD

A CAD of \mathbb{R}^n is a **partition** of \mathbb{R}^n , where

- ▶ all the cells are **cylindrically** arranged, that is for all $1 \leq j < n$ the projections on the first j coordinates (y_1, \dots, y_j) of any two cells are either identical or disjoint.
- ▶ each cell is a **connected semi-algebraic** subset, called a region

For $F_n \subset \mathbb{Q}[y_1, \dots, y_n]$, a CAD of \mathbb{R}^n is **F_n -invariant** if above each region of it, the sign of each $f \in F_n$ is constant.

Our Method

F_n : a set of polynomials of $\mathbb{Q}[y_1, \dots, y_n]$.

Initial Partition: we decompose \mathbb{C}^n into disjoint constructible sets C_1, \dots, C_e such that for each $f \in F_n$, either f is identically zero in C_j or f vanishes at no points of C_j .

Make Cylindrical: we transform the initial partition and obtain another partition of \mathbb{C}^n into disjoint constructible sets such that this second decomposition is **cylindrical**.

Make Semi-Algebraic: from the previous decomposition we produce an F_n -invariant CAD of \mathbb{R}^n via real root isolation of zero-dimensional regular chains.

The Three Phases

$$F_n \subset \mathbb{Q}[y_1, \dots, y_n]$$



Initial Partition



\mathcal{C} : a partition of \mathbb{C}^n into constructible sets



Make Cylindrical



\mathcal{D} : a cylindrically arranged partition of \mathbb{C}^n into constructible sets



Make SemiAlgebraic



An F_n -invariant CAD of \mathbb{R}^n

Representation of Constructible Sets

A pair $R = [T, h]$ is called a **regular system** if T is a regular chain, and h is a polynomial which is regular w.r.t $\text{sat}(T)$.

Theorem (CGLMP, CASC2007)

Every constructible set can be written as a finite union of the zero sets of regular systems.

The constructible set

$$\begin{cases} x(1+y) - s = 0 \\ y(1+x) - s = 0 \\ x + y - 1 \neq 0 \end{cases} \quad (1)$$

can be represented by two regular systems

$$R_1 : \left| \begin{array}{l} T_1 = \begin{cases} (y+1)x - s \\ y^2 + y - s \end{cases} \\ h_1 = y - 2s + 1 \end{array} \right. \quad R_2 : \left| \begin{array}{l} T_2 = \begin{cases} x + 1 \\ y + 1 \\ s \end{cases} \\ h_2 = 1 \end{array} \right.$$

Initial Partition

Let $F_n = \{f_1, \dots, f_s\}$ be a finite subset of $\mathbb{Q}[y_1 < \dots < y_n]$. We compute an **intersection free basis** of the $s + 1$ sets $f_1 = 0, \dots, f_s = 0$ and $f_1 \cdots f_s \neq 0$, where each element is represented a regular system and their sets from a partition of \mathbb{C}^n .

Consider the parametric parabola $p = ax^2 + bx + c$, where $x > c > b > a$. **InitialPartition** decomposes \mathbb{C}^4 into four pairwise disjoint sets, each of which is the zero set of a regular system.

$$r_1 := \begin{cases} c = 0 \\ b = 0 \\ a = 0 \end{cases}, \quad r_2 := \begin{cases} bx + c = 0 \\ b \neq 0 \\ a = 0 \end{cases},$$

$$r_3 := \begin{cases} ax^2 + bx + c = 0 \\ a \neq 0 \end{cases}, \quad r_4 := \{ ax^2 + bx + c \neq 0 \}.$$

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Separate Zeros

Let $rs = [T, h]$ be a regular system of $\mathbb{Q}[y_1 < \cdots < y_n]$. We see y_1, \dots, y_{n-1} as parameters, denoted by \mathbf{u} , and regard rs as a parametric system in \mathbf{u} , that we solve via **comprehensive triangular decomposition**.

As a result, we obtain a partition of the projection onto the \mathbf{u} -space of $Zero(rs)$ such that, above each cell R of the partition, $Zero(rs)$ equals the union of the zero sets of some polynomials $p_1, \dots, p_r \in \mathbb{R}[y_1, \dots, y_n]$, where

- ▶ the initial of each p_j does not vanish on R ,
- ▶ the p_j 's are squarefree and pairwise coprime at any point of R .

For the regular system

$$r_3 := \begin{cases} ax^2 + bx + c = 0 \\ a \neq 0 \end{cases}$$

Calling `SeparateZeros(r_3)` will get

$$\begin{aligned} \{a(4ac - b^2) \neq 0\} &\rightarrow \{ax^2 + bx + c\} \\ \{4ac - b^2 = 0, a \neq 0\} &\rightarrow \{2ax + b\} \end{aligned}$$

Make Cylindrical

By calling `SeparateZeros` recursively, `MakeCylindrical` produces a **cylindrical decomposition** of \mathbb{C}^n , defined inductively as follows.

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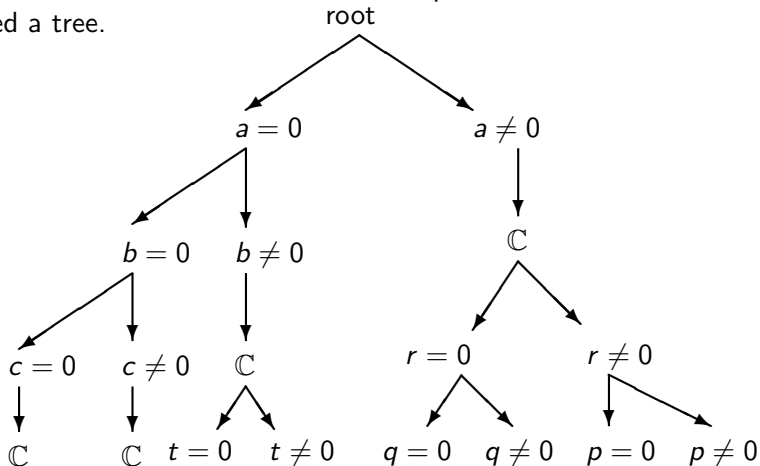
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 - ▶ the initial of each p_j does not vanish on D_i and,
 - ▶ the p_j 's are squarefree and pairwise coprime at all $\mathbf{u} \in D_i$,
 - ▶ $D_i \times (p_1 = 0), \dots, D_i \times (p_r = 0)$ and $D_i \times (p_1 \cdots p_r \neq 0)$ are in \mathcal{D} .

The algorithm **MakeCylindrical** takes r_1, r_2, r_3 and r_4 as input and outputs a cylindrical decomposition of \mathbb{C}^4 . Let $t = bx + c$, $q = 2ax + b$, and $r = 4ac - b^2$, the decomposition can be described a tree.



The Three Phases

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Make SemiAlgebraic



An F_n -invariant CAD of \mathbb{R}^n

Make SemiAlgebraic (I)

Theorem (Collins)

*Let $p \in \mathbb{R}[y_1 < \cdots < y_n]$ and R be a region of \mathbb{R}^{n-1} . If $\text{init}(p)$ does not vanish on R and the number of distinct complex roots of p is invariant on R , then p is **delineable** on R , that is, $V(p)$ is the union of finitely many disjoint graphs of continuous functions over R .*

Make SemiAlgebraic (II)

Corollary

Let $\{p_1, \dots, p_r\} \subset \mathbb{R}[y_1 < \dots < y_n]$ and let R be a region of \mathbb{R}^{n-1} . Assume that for all $\alpha \in R$:

- ▶ each $\text{init}(p_j)$ does not vanish at α ;
- ▶ all $p_j(\alpha, y_n)$, as polynomials of $\mathbb{R}[y_n]$, are squarefree and coprime.

Then each p_j is delineable on R and any two sections of the cylinder over R , given by different p_i and p_j , are disjoint.

By Collins' theorem and its corollary, one derives a CAD of \mathbb{R}^n from a cylindrical decomposition of \mathbb{C}^n , by means real root isolation of zero-dimensional regular chains.

Maple Demo

Special thanks to James H. Davenport and John May for the piecewise construction.

Comparing with Collins' Algorithm

Consider the parametric parabola $p = ax^2 + bx + c$, where $x > c > b > a$.

- ▶ Our algorithm produces a p -invariant CAD of \mathbb{R}^4 with 27 cells, which is minimal.
- ▶ By Collins-Hong or McCallum projection operator, one produces the following polynomials during the projection phase:

$$ax^2 + bx + c, b^2 - 4ac, c, b, a.$$

In the lifting phase, one then obtains a CAD of \mathbb{R}^4 with 115 cells (Brown 01)!

- ▶ If Brown-McCallum projection operator is applied, one could also obtain a CAD of \mathbb{R}^4 with 27 cells (Brown 01). However, this projection operator may fail in some (rare) cases.

| Sys | InitialPartition | MakeCylindrical | MakeSemiAlgebraic | Total | $N_{\mathbb{R}}$ |
|-----|------------------|-----------------|-------------------|----------|------------------|
| 1 | 0.024 | 0.096 | 0.024 | 0.144 | 27 |
| 2 | 1.184 | 2.856 | 1.048 | 5.088 | 895 |
| 3 | 0.004 | 7.512 | 0.704 | 8.220 | 233 |
| 4 | 0.264 | 1.368 | 1.080 | 2.716 | 421 |
| 5 | 0.016 | 0.052 | 0.116 | 0.184 | 55 |
| 6 | 0.108 | 0.156 | 0.120 | 0.384 | 41 |
| 7 | 2.704 | 3.600 | 1.360 | 7.664 | 893 |
| 8 | 0.380 | 1.608 | 1.196 | 3.184 | 365 |
| 9 | 0.288 | 0.532 | 0.264 | 1.084 | 209 |
| 10 | 5.668 | 48.079 | 18.833 | 72.640 | 3677 |
| 11 | 0.252 | 1.192 | 0.620 | 2.068 | 563 |
| 12 | 2.664 | 135.028 | 88.142 | 225.862 | 20143 |
| 13 | 10.576 | 35.846 | 6.905 | 53.335 | 4949 |
| 14 | 5.728 | 71.760 | 2520.354 | 2597.878 | 27547 |
| 15 | 690.731 | 2513.817 | 299.250 | 3503.954 | 66675 |
| 16 | 895.435 | 2064.469 | - | - | - |
| 17 | 0.052 | - | - | - | - |
| 18 | - | - | - | - | - |

Table 1 Timing (s) and number of cells for CAD

Observation

- ▶ For most examples the steps of the algorithm dedicated to computations in complex space, where GCDs of polynomials modulo regular chains are computed intensively, dominate the step taking place in the real space.
- ▶ The data suggests that the modular methods and efficient implementation techniques being developed in `RegularChains` library have a large potential for improving our current implementation.

Conclusion

- ▶ We have introduced an intermediate concept, **cylindrical decomposition of the complex space**, from which a CAD of \mathbb{R}^n can easily be extracted.
- ▶ W.r.t Collins-Hong projection operator, even for simple examples, our approaches tends to produce **much less cells** due to its case discussion feature.
- ▶ W.r.t Brown-McCallum projection operator, it **can always generate** a CAD while the Brown-McCallum projection operator may fail (rarely).

1. Parametric parabola: $\{ax^2 + bx + c\}, x > c > b > a.$
2. Whitney umbrella: $\{x - uv, y - v, z - u^2\}, v > u > z > y > x.$
3. Quartic: $\{x^4 + px^2 + qx + r\}, x > p > q > r.$
4. Sphere-Catastrophe: $\{z^2 + y^2 + x^2 - 1, z^3 + xz + y\}, x > y > z.$
5. Tacnode curve: $\{y^4 - 2y^3 + y^2 - 3x^2y + 2x^4\}, y > x.$
6. Arnón-84-2: $\{144y^2 + 96x^2y + 9x^4 + 105x^2 + 70x - 98, xy^2 + 6xy + x^3 + 9x\}, y > x.$
7. A real implicitization problem:
 $\{x - uv, y - uv^2, z - u^2\}, v > u > z > y > x.$
8. Ball-circular-cylinder:
 $\{x^2 + y^2 + z^2 - 1, x^2 + (y + z - 2)^2 - 1\}, z > y > x.$
9. Termination of term rewrite system
 $\{x - r, y - r, x^2(1 + 2y)^2 - y^2(1 + 2x^2)\}, r > x > y.$
10. Collins and Johnson: $\{3a^2r + 3b^2 - 2ar - a^2 - b^2, 3a^2r + 3b^2r - 4ar + r - 2a^2 - 2b^2 + 2a, a - 1/2, b, r, r - 1\}, r > a > b.$
11. Range of lower bounds
 $\{a, az^2 + bz + c, ax^2 + bx + c - y\}, z > c > b > a > x > y.$

12. X-axis ellipse problem: $\{b^2(x - c)^2 + a^2y^2 - a^2b^2, x^2 + y^2 - 1\}, y > x > b > c > a.$

13. Davenport and Heintz

$\{a - d, b - c, a - c, b - 1, a^2 - b\}, a > b > c > d.$

14. Hong-90

$\{r + s + t, rs + st + tr - a, rst - b\}, t > s > r > b > a.$

15. Solotareff-3

$\{r, r - 1, u + 1, u - v, v - 1, 3u^2 + 2ru - a, 3v^2 + 2rv - a, u^3 + ru^2 - au + a - r - 1, v^3 + rv^2 - av - 2b - a + r + 1\}, b > u > v > r > a.$

16. Collision problem

$\{\frac{17}{16}t - 6, \frac{17}{16}t - 10, x - \frac{17}{16}t + 1, x - \frac{17}{16}t - 1, y - \frac{17}{16}t + 9, y - \frac{17}{16}t + 7, (x - t)^2 + y^2 - 1\}, t > x > y.$

17. McCallum trivariate random polynomial

$\{(y - 1)z^4 + xz^3 + x(1 - y)z^2 + (y - x - 1)z + y\}, z > y > x.$

18. Ellipse problem

$\{b^2(x - c)^2 + a^2(y - d)^2 - a^2b^2, a, b, x^2 + y^2 - 1\}, y > x > d > c > b > a.$