



# Change of order for regular chains in positive dimension

Xavier Dahan<sup>\*</sup>, Xin Jin<sup>†</sup>, Marc Moreno Maza<sup>†</sup>, and Éric Schost<sup>†</sup>

<sup>\*</sup>: LIX, École polytechnique, 91128 Palaiseau, France.

<sup>†</sup>: ORCCA, University of Western Ontario, London, Canada.

ISSAC, Genova, July 9-12th, 2006.



## Motivation and examples

We describe an algorithm for changing the order of **regular chains**<sup>\*</sup> in positive dimension.

**Example: invariant theory.** Polynomials  $P(X, Y)$  invariant under  $(X, Y) \mapsto (-X, -Y)$  can be written in terms of

$$S = X^2, T = Y^2, U = XY.$$

To rewrite an invariant polynomial, it helps to obtain the expression of  $X$  and  $Y$  in term of  $S, T, U$ . This is done by **changing the order** in the input system

$$\begin{cases} \mathbf{S} - X^2 \\ \mathbf{T} - Y^2 \\ \mathbf{U} - XY \end{cases} \quad \text{for } \mathbf{S} > \mathbf{T} > \mathbf{U} > \mathbf{X} > \mathbf{Y}$$

$$\text{to } \begin{cases} \mathbf{S} \mathbf{Y} - \mathbf{U} \mathbf{X} \\ \mathbf{X}^2 - \mathbf{S} \\ \mathbf{U}^2 - \mathbf{S} \mathbf{T} \end{cases} \quad \text{for } \mathbf{Y} > \mathbf{X} > \mathbf{U} > \mathbf{S} > \mathbf{T}.$$

This is also an example of an **implicitization** problem.

**Main result.** Let  $C$  be a regular chain, whose **saturated ideal**<sup>\*</sup> is prime.

We give a probabilistic algorithm, of complexity polynomial in the size of input (**degree, complexity of evaluation**) and output (**number of monomials**), and in the degree of the **quasi-component**<sup>\*</sup> of  $C$ .

Let  $d$  be the maximum degree in  $C$ , let  $n$  be the number of variables. If all random choices are made uniformly in a finite set  $\Gamma$ , the probability of failure is at most  $\frac{2n(3d^n+n^2)d^{2n}}{|\Gamma|}$ .

**Specificities.** We use a small set of **well-defined subroutines**.

- change of order in dimension zero;
- Newton-Hensel lifting.

### Definition (regular chain, quasi-component)

Let  $X$  be  $n$  ordered variables.

Let  $C = C_1, \dots, C_s$  be in  $k[X]$ , with main variables  $X_{\ell_1} < \dots < X_{\ell_s}$ .

- the **initial**  $h_i$  is the leading coefficient of  $C_i$  in  $X_{\ell_i}$ ;
- the  $i$ th **saturated ideal** is  $\text{Sat}_i = \langle C_1, \dots, C_i \rangle : (h_1 \cdots h_i)^\infty$ ;
- $C$  is a **regular chain** if  $h_i$  is regular mod  $\text{Sat}_i$  for all  $i$ ;
- the **quasi-component** is the zero-set of  $\text{Sat}_n$ .

## Basic setup

Let  $V$  be an irreducible variety of dimension  $r$ , defined by polynomials in  $k[X] = k[X_1, \dots, X_n]$ .

**Free / algebraic variables.** A set  $Y$  of  $r$  variables  $Y_1, \dots, Y_r$  is **free** if the image

$$(x_1, \dots, x_n) \in V \mapsto (y_1, \dots, y_r)$$

is dense. Then, the generic points of  $V$  can be described by a regular chain

$$\begin{cases} T_s(Y, Z_1, \dots, Z_s) \\ \vdots \\ T_1(Y, Z_1) \end{cases} \quad \text{with } Z = (Z_1, \dots, Z_s) = X - Y$$

(**algebraic variables**)

**Data structure.** If  $Y$  are free variables, we represent  $V$  by

$$y = (y_1, \dots, y_r) \in k^r, \quad \begin{cases} T_s(y, Z_1, \dots, Z_s) \\ \vdots \\ T_1(y, Z_1) \end{cases}$$

**Specialize and lift** paradigm: intermediate computations are done in **dimension 0 and 1**.

**Exchange property.** Let  $Y$  and  $Y'$  be sets of free variables, and let  $v \in Y - Y'$ . Then, there exists  $w \in Y' - Y$  such that  $Y - v + w$  is free.

This means that the sets of free variables form a **matroid**<sup>\*</sup>, the **coordinate matroid** of  $V$ .

Algorithmically, given  $T_1(y, Z_1), \dots, T_s(y, Z_1, \dots, Z_s)$ , we use a subroutine called **Swap**:

1. change of order in  $T$  to put  $w$  as last variable; [dim. 0]
2. lift  $v$  in  $T$ ; [dim. 1]
3. specialize  $w$  at a random value. [dim. 0]

### Definition (matroid)

Let  $X$  be a finite set.

A matroid over  $X$  is a collection  $B(M)$  of subsets of  $X$  with the same cardinality  $r$  and satisfying the **exchange property**:

For all  $Y$  and  $Y'$  in  $B(M)$ , for every  $v \in Y - Y'$  there exists  $w \in Y' - Y$  such that  $Y - v + w$  is in  $B(M)$ .

## Algorithm

**Input.**

- A **regular chain**  $T_{\text{in}}$ , whose saturated ideal  $I$  is prime.
- A **target order**  $\succ$  on the variables.

**Output.**

- A **regular chain**  $T_{\text{out}}$  for the order  $\succ$ , with saturated ideal  $I$ .

**Step 1.** Determine a sequence of **exchanges**.

**Prop.** For a generic  $\alpha \in V$ , the coordinates matroids of (i) the variety  $V$  and (ii) the tangent space  $T_\alpha V$  coincide.

**Prop.** The algebraic variables in  $T_{\text{out}}$  are the maximal set of algebraic variables for a suitable lexicographic order.

Then, using **linear algebra** in the Jacobian matrix at  $\alpha$ , we can find by a greedy algorithm a sequence of exchanges:

$$Y_0 = Y_{\text{in}} \quad (\text{free variables of } T_{\text{in}})$$

$$\mapsto Y_1 = Y_0 - v_0 + w_0$$

...

$$\mapsto Y_N = Y_{N-1} - v_{N-1} + w_{N-1} \quad (\text{free variables of } T_{\text{out}})$$

**Step 2.** Apply the exchanges to the triangular sets.

$$T_0 \quad (\text{specialization of } T_{\text{in}})$$

$$\mapsto T_1 = \text{Swap}(T_0, v_0, w_0)$$

...

$$\mapsto T_N = \text{Swap}(T_{N-1}, v_{N-1}, w_{N-1}) \quad (\text{specialization of } T_{\text{out}})$$

**Step 3.** Lift all free variables in  $T_N$ .

**Implementation.** This algorithm is implemented in the **RegularChains** Maple package.

### References

**Regular chains** Aubry, Kalkbrener, Lazard, Moreno Maza, Lemaire-Moreno Maza-Xie (**RegularChains**), ...

**Lifting / specialization** Giusti, Heintz, Lecerf, Pardo *et al.*

**Change of order** Faugère-Gianni-Lazard-Mora, Collart-Kalkbrener-Mall (Gröbner walk), Boulier-Lemaire-Moreno Maza (Pardi)

**Implicitization** Buzé, Chardin, Cox, D'Andrea, Jouanolou, Khetan, ...