Polynomials over Power Series and their Applications to Symbolic Analysis

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Does the parametrization reach all points of the surface? (1/8)



Figure: Steiner's Roman surface

https://upload.wikimedia.org/wikipedia/commons/e/ea/Steiner%27s_Roman_Surface.gif An implicit formula of Steiner's Roman surface S is f = 0, where:

$$f := 4 x^{4} - 8 yx^{3} + 9 x^{2}y^{2} - 8 yzx^{2} - 5 y^{3}x + 8 y^{2}zx + y^{4}$$

$$-2 y^{3}z + 3 y^{2}z^{2} - 2 yz^{3} + z^{4} - 8 yx^{2} + 8 zx^{2} + 8 y^{2}x$$
(1)
$$-8 xyz - 2 y^{3} + 2 y^{2}z - 2 yz^{2} + 4 x^{2} - 4 yx + y^{2}.$$

Does the parametrization reach all points of the surface? (2/8)

• With $q(s,t) := s^2 + t^2 + s - t + 1$, consider also the following map

$$\vec{r}: \quad \mathbb{R}^2 \quad \to \qquad \mathbb{R}^3 \\ (s,t) \quad \mapsto \quad \left(\frac{s^2}{q(s,t)} \ , \ \frac{s^2+t^2}{q(s,t)} \ , \ \frac{s^2+s\,t+s+t}{q(s,t)}\right), \tag{2}$$

- Do we have $\text{Image}(\vec{r}) = S$?
- A preliminary question is whether q(s,t) vanishes or not.

>
$$R \coloneqq PolynomialRing([s, t, x, y, z]): q \coloneqq s^2 + t^2 + s - t + 1:$$

RealTriangularize([q], R);
[]

Figure: RegularChains:-RealTriangularize proves q(s,t) has no real points.

Does the parametrization reach all points of the surface? (3/8)

Let us verify that the image of the map \vec{r} is contained in the surface S.

> $f := 4 \cdot x^4 - 8 \cdot y \cdot x^3 + 9 \cdot x^2 \cdot y^2 - 8 \cdot y \cdot z \cdot x^2 - 5 \cdot y^3 \cdot x + 8$ $y^{2} \cdot z \cdot x + y^{4} - 2 \cdot y^{3} \cdot z + 3 \cdot y^{2} \cdot z^{2} - 2 \cdot y \cdot z^{3} + z^{4} - 8 \cdot y \cdot x^{2} + 8 \cdot z \cdot x^{2} + 8 \cdot y^{2} \cdot x - 8 \cdot x \cdot y \cdot z - 2$ $v^{3} + 2v^{2}v^{2} - 2v^{2}v^{2} + 4v^{2} + 4v^{2} + 4v^{2} + 4v^{2}$ $R \coloneqq PolynomialRing([s, t, x, y, z])$: dec1 := Triangularize([f], R); S := GeneralConstruct(dec1[1], map(Initial, Map(Initial)))Equations (dec1[1], R), R), R); $dec1 := [regular_chain]$ $S := constructible_set$ > $q := s^2 + t^2 + s - t + 1$: $F := [a \cdot x - s^2, a \cdot y - (s^2 + t^2), a \cdot z - (s^2 + s \cdot t + s + t)];$ dec2 = Triangularize(F, R); ImageR = GeneralConstruct(dec2[1], map(Initial, F, R), R); $dec2 := [reaular_chain]$ ImaaeR := constructible_set > LM1 ≔ Difference(ImageR, S, R); IsEmpty(LM1, R); LM1 := constructible settrue

Figure: The command Difference computes the points in the image of \vec{r} that do not belong to surface S, which is empty.

Does the parametrization reach all points of the surface? (4/8)

- Disproving $\text{Image}(\vec{r}) = S$ can be done by specialization
- Computing $\operatorname{Image}(\vec{r}) \cap \{y = 1\}$ yields

$$2x^2 + 2xz + z^2 - 3x - 2z + 1 = 0$$

• While computing $S \cap \{y = 1\}$ brings more:

 $(2x^{2} - 2xz + z^{2} - x)(2x^{2} + 2xz + z^{2} - 3x - 2z + 1) = 0$



Does the parametrization reach all points of the surface? (5/8)

$$\begin{array}{l} > R \coloneqq PolynomialRing([s, t, x, y, z]):\\ q \coloneqq s^{\lambda}2 + t^{\lambda}2 + s - t + 1:\\ F \coloneqq [x^{\lambda}q - s^{\lambda}Q, y^{*}q - (s^{\lambda}2 + t^{\lambda}2), z^{*}q - (s^{\lambda}2 + s^{*}t + s + t)]:\\ dec2 \coloneqq Projection([op(F), y - 1], [], [], [], 3, R): Display(\%, R)\\ \left[\left(\begin{array}{c} 4 x + 2 z - 3 = 0\\ y - 1 = 0\\ 4 z^{2} - 4 z - 1 = 0 \end{array}\right) \left(\begin{array}{c} x = 0\\ y - 1 = 0\\ z - 1 = 0\end{array}\right) \left(\begin{array}{c} 2 x^{2} + (2 z - 3) x + z^{2} - 2 z + 1 = 0\\ y - 1 = 0\\ 4 z^{2} - 4 z < 1 \text{ and } z - 1 \neq 0 \end{array}\right) \end{array}\right)$$

$$\begin{cases} > f \coloneqq 4x^4 - 8yx^3 + (9y^2 + (-8z - 8)y + 8z + 4)x^2 + (-5y^3 + (8z + 8)y^2 + (-8z - 4)y)x \\ + y^4 + y^3 (-2z - 2) + (3z^2 + 2z + 1)y^2 + (-2z^3 - 2z^2)y + z^4 \\ R \vDash PolynomialRing([s, t, x, y, z]) \\ decl \coloneqq RealTriangularize([f, y - 1], R) : Display(decl, R); \\ [2x^2 + (2z - 3)x + z^2 - 2z + 1 = 0 \\ y - 1 = 0 \\ 4z^2 - 4z < 1 \\ \end{cases} \begin{pmatrix} 4x + 2z - 3 = 0 \\ y - 1 = 0 \\ 4z^2 - 4z - 1 = 0 \\ \end{cases}$$

Does the parametrization reach all points of the surface? (6/8)

Difference(dec1, dec2, R): Display(%, R);			
x - 1 = 0	x = 0	2x - 1 = 0	4x - 2z - 1 = 0
y - 1 = 0 ,	y-1=0 ,	$y-1=0 \qquad , \qquad$	y - 1 = 0
z - 1 = 0	<i>z</i> = 0	z - 1 = 0	$4 z^2 - 4 z - 1 = 0$
$2x^{2} + (-1 - 2z)x + z^{2} = 0$			
y - 1 = 0			
$4 z^2 - 4 z < 1$ and $z \neq 0$ and $z - 1 \neq 0$ and $2 z - 1 \neq 0$			

Figure: The points on Steiner surface S and the plane y = 1 which do not belong to the intersection of the image of the parametrization \vec{r} and the plane y = 1.

Observe that these calculations are done over the **reals**!

The next question

- Therefore, $\text{Image}(\vec{r}) = S \text{ does not hold}!$
- 2 Next question: can we recover from S what $Image(\vec{r})$ is missing?
- \bullet if the missing point are $\overline{\text{Image}(\vec{r})} \setminus \text{Image}(\vec{r})$, then the answer is yes.

The closure of a constructible set

- Denote by $\overline{\text{Image}(\vec{r})}$ the closure of $\text{Image}(\vec{r})$ in the Euclidean topology (over \mathbb{C}).
- **2** Thanks to a theorem of David Mumford, $\overline{\text{Image}(\vec{r})}$ is also the closure of $\text{Image}(\vec{r})$ in Zariski topology.
- Thus Image(r) is the intersection of all algebraic sets containing Image(r).
- By the way, Gröbner basis techniques can capture Zariski closures over algebraically closed fields.

Does the parametrization reach all points of the surface? (8/8)

```
> q := s^2 + t^2 + s - t + 1:

R := [x^*q - s^2, y^*q - (s^2 + t^2), z^*q - (s^2 + s^* t + s + t)]:

with (PolynomialIdeals):

sat := Saturate((op(R)), q):

closure_of_Image_of_r := EliminationIdeal(sat, {x, y, z});

closure_of_Image_of_r := (4 x^4 - 8 x^3 y + 9 x^2 y^2 - 8 x^2 y z - 5 x y^3 + 8 x y^2 z + y^4 - 2 y^3 z + 3 y^2 z^2 - 2 y z^3 + z^4 - 8 x^2 y + 8 x^2 z + 8 x y^2 - 8 x y z - 2 y^3 + 2 y^2 z - 2 y z^2 + 4 x^2 - 4 x y + y^2)
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Figure: Closure of Image(\vec{r}).

- We retrieve the polynomial defining the implicit representation of ${\cal S}$
- According to the so-called *Elimination Theorem* (see the book *Ideals*, varieties and Algorithms) the algebraic set of the elimination ideal $\mathcal{I} \subset \mathbb{K}[x_1 < \cdots < x_n]$ w.r.t. x_1, \ldots, x_k (for some $1 \leq k < n$) is equal to the Zariski closure of the projection of $V(\mathcal{I})$ onto x_1, \ldots, x_k .

- Computing Zariski closures of constructible sets (that is, systems of polynomial equations and inequation) and semi-algebraic sets (that is, systems of polynomial equations and inequalities) appear naturally in practice: reachable sets, projection of constructible sets and semi-algebraic sets.
- Gröbner basis techniques can deal with the case of constructible sets.
- We are mainly interested here with the real case, that is, semi-algebraic sets .

Topological closure and limit points

Let (X, τ) be a topological space and $S \subseteq X$ be a subset.

Topological closure

The <u>closure</u> of S, denoted \overline{S} , is the intersection of all closed sets containing S.

Limit point

- A point $p \in X$ is a *limit point* of S if every neighbourhood of p contains at least one point of S different from p itself.
- The limit points of S which do not belong to S are called non-trivial, denoted by $\lim(S)$.

Properties

• If X is a metric space, the point p is a limit point of S if and only if there exists a sequence $(x_n, n \in \mathbb{N})$ of points of $S \setminus \{p\}$ such that $\lim_{n \to \infty} x_n = p$.

• We have
$$\overline{S} = S \cup \lim(S)$$
.

Zariski topology and the Euclidean topology

The relation between the two topologies

- With $\mathbb{K} = \mathbb{C}$, the affine space \mathbb{A}^s is endowed with both topologies.
- The basic open sets of the Euclidean topology are the open balls.
- The basic open sets of Zariski topology are the complements of hypersurfaces.
- Thus, a Zariski closed (resp. open) set is closed (resp. open) in the Euclidean topology on \mathbb{A}^s .
- That is, Zariski topology is coarser than the Euclidean topology.

The relation between the two closures (D. Mumford)

- Let $V \subseteq \mathbb{A}^s$ be an irreducible affine variety.
- Let $U \subseteq V$ be nonempty and open in Zariski topology induced on V.

Then, \boldsymbol{U} has the same closure in both topologies. In fact, we have

$$V = \overline{U}^Z = \overline{U}^E.$$

Limit points: a first example

- Let S be the zero-set of a polynomial system and \overline{S} be the topological closure S in the Euclidean topology.
- It can be proved that the set-theoretic difference $\overline{S}\setminus S$ can be obtained via a *limit computation process* illustrated below

Consider S below together with a Puiseux series expansion around z = 0:

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$$S := \begin{cases} z x - y^2 = 0\\ y^5 - z^4 = 0\\ z \neq 0 \end{cases} \quad \text{and} \quad \begin{cases} x = \frac{t^{6/3}}{t}\\ y = t^{4/5}\\ z = t\\ t \neq 0 \end{cases}$$

Then we have:

$$\lim_{t \to 0} \begin{pmatrix} x(t) \\ y(t) \\ z(t) \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix} \text{ and } \overline{S} \setminus S = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$$

Limit points: a second example

Consider S below together with a Puiseux series expansion around z = 0:

$$S := \begin{cases} z x - y^2 = 0\\ y^5 - z^2 = 0\\ z \neq 0 \end{cases} \text{ and } \begin{cases} x = t^{-1/5}\\ y = t^{2/5}\\ z = t\\ t \neq 0 \end{cases}$$

Then we have:

$$\lim_{t \to 0} \begin{pmatrix} x(t) \\ y(t) \\ z(t) \end{pmatrix} = \begin{pmatrix} \pm \infty \\ 0 \\ 0 \end{pmatrix} \text{ and } \overline{S} \setminus S = \emptyset$$

Regular chains in a nutshell

- Regular chains generalize the concept of *triangular system* from linear algebra to polynomial algebra.
- Thus, they are polynomial systems with a triangular shape and additional algebraic properties which support a back substitution process.
- Every (non-constant) bivariate polynomial forms a regular chain.

The solutions of a regular chain

- Like Gröbner bases, regular chains can be used to compute and describe the solutions of polynomial systems over algebraically closed fields, say C.
- \bullet Regular chains can also be used to solve over real closed fields, say $\mathbb R$ but also Puiseux series.

The Puiseux series solutions of a regular chain (2/2)

Limit points: yet another example

 $\begin{aligned} R &:= PolynomialRing([x, y, z]): rc := Chain([-y^3 + y^2 + z^5, z^4 * x + y^3 - y^2], Empty(R), R): \\ Display(rc, R); \end{aligned}$

$$z4 x + y3 - y2 = 0$$
$$-y3 + y2 + z5 = 0$$
$$z4 \neq 0$$

RegularChainBranches(rc, R, [z]);

 $\begin{bmatrix} z = T^2, y = \frac{1}{2} T^5 (-T^5 + 2 \operatorname{RootOf}(-Z^2 + 1)), x = -\frac{1}{8} T^2 (-T^{20} + 6 T^{15} \operatorname{RootOf}(-Z^2 + 1) + 10 T^{10} + 8) \end{bmatrix}, \begin{bmatrix} z = T^2, y = -\frac{1}{2} T^5 (T^5 + 2 \operatorname{RootOf}(-Z^2 + 1)), x = \frac{1}{8} T^2 (T^{20} + 6 T^{15} \operatorname{RootOf}(-Z^2 + 1) - 10 T^{10} - 8) \end{bmatrix}, \begin{bmatrix} z = T, y = T^5 + 1, x = -T (T^{10} + 2 T^5 + 1) \end{bmatrix}$

lp := LimitPoints(rc, R) : Display(lp, R);

$$\begin{cases} x = 0 \\ y = 0 \\ z = 0 \end{cases}, \begin{cases} x = 0 \\ y - 1 = 0 \\ z = 0 \end{cases}$$

Figure: Computation of (non-trivial) limit points with the RegularChains library

Limit points: statement of our quest

• Let
$$R := \{t_2(x_1, x_2), \dots, t_n(x_1, \dots, x_n)\}$$

- We regard t_i as a univariate polynomial w.r.t. x_i , for i = 2, ..., n:
- We denote by h_i the leading coefficient (also called initial) of t_i w.r.t. x_i , and assume that h_i depends on x_1 only; hence we have

 $t_i = h_i(x_1)x_i^{d_i} + c_{d_i-1}(x_1, \dots, x_{i-1})x_i^{d_i-1} + \dots + c_0(x_1, \dots, x_{i-1})$ • Consider the system

$$W(R) := \begin{cases} t_n(x_1, \dots, x_n) = 0 \\ \vdots \\ t_2(x_1, x_2) = 0 \\ (h_2 \cdots h_n)(x_1) \neq 0 \end{cases}$$

Main Goal

- Where do the points of W(R) go when x_1 approaches a root of $h_2 \cdots h_n$?
- In other words, we want to compute the points which belong to the topological closure of W(R) but to W(R) itself.

Limit points: yet again another example

> R := PolynomialRing([x, y, z]): $rc \coloneqq Chain([y^{(3)}-2^*y^{(3)}+y^{(2)}+z^{(5)},z^{(4)}*x+y^{(3)}-y^{(2)}], Empty(R), R): Display(rc, R);$ br := RegularChainBranches(rc. R. [z], coefficient = complex); $\begin{cases} z^4 x + y^3 - y^2 = 0 \\ -y^3 + y^2 + z^5 = 0 \\ z^4 \neq 0 \end{cases}$ $br := \left[\left[z = T^2, y = \frac{1}{2} T^5 \left(-T^5 + 2 \operatorname{RootOf}(_2 Z^2 + 1) \right), x = -\frac{1}{8} T^2 \left(-T^{20} + 6 T^{15} \operatorname{RootOf}(_2 Z^2 + 1) + 10 T^{10} + 8 \right) \right] \right]$ $\left[z = T^{2}, y = -\frac{1}{2}T^{5}\left(T^{5} + 2 \operatorname{RootOf}(_{Z}^{2} + 1)\right), x = \frac{1}{8}T^{2}\left(T^{20} + 6 T^{15}\operatorname{RootOf}(_{Z}^{2} + 1) - 10 T^{10} - 8\right)\right], \left[z = T^{2}, y = -\frac{1}{2}T^{5}\left(T^{5} + 2 \operatorname{RootOf}(_{Z}^{2} + 1)\right), x = \frac{1}{8}T^{2}\left(T^{20} + 6 T^{15}\operatorname{RootOf}(_{Z}^{2} + 1) - 10 T^{10} - 8\right)\right], \left[z = T^{2}, y = -\frac{1}{2}T^{5}\left(T^{5} + 2 \operatorname{RootOf}(_{Z}^{2} + 1)\right), x = \frac{1}{8}T^{2}\left(T^{20} + 6 T^{15}\operatorname{RootOf}(_{Z}^{2} + 1)\right) - 10T^{10} - 8\right)\right], \left[z = T^{2}, y = -\frac{1}{2}T^{5}\left(T^{5} + 2 \operatorname{RootOf}(_{Z}^{2} + 1)\right), x = \frac{1}{8}T^{2}\left(T^{20} + 6 T^{15}\operatorname{RootOf}(_{Z}^{2} + 1)\right) - 10T^{10} - 8\right)\right]$ $= T, y = T^{5} + 1, x = -T (T^{10} + 2 T^{5} + 1)]$ > br := RegularChainBranches(rc, R, [z], coefficient = real); $br := \left[\left[z = T, y = T^5 + 1, x = -T \left(T^{10} + 2 T^5 + 1 \right) \right] \right]$

Figure: The command RegularChainBranches computes a parametrization for the complex and real paths of the quasi-component defined by rc. When coefficient argument is set as real, then the command RegularChainBranches computes the real branches.

Application 1: limit of multivariate rational functions



Figure: On the left: the surface defined by $q := \frac{x^4+3x^2y-x^2-y^2}{x^2+y^2} = z$ around the origin. On the right: the three paths of discriminant variety of q going through the point (0,0,-1).

Application 2: tangent cone computations



Figure: The tangent cone of the "fish" given by $f := y^2 - x^2 (x + 4) = 0$ at the origin consists of two tangent lines: y = 2x and y = -2x.

Application 3: computing intersection multiplicities

> $F := [(x^2 + y^2)^2 + 3x^2y - y^3, (x^2 + y^2)^3 - 4x^2y^2]$: > plots[implicitplot](Fs, x = -2..2, y = -2..2): (3)> R := PolynomialRing([x, y], 101) :> TriangularizeWithMultiplicity(F, R); $\begin{bmatrix} \begin{bmatrix} 1, \begin{cases} x-1=0\\ y+14=0 \end{bmatrix} \end{bmatrix}, \begin{bmatrix} \begin{bmatrix} 1, \begin{cases} x+1=0\\ y+14=0 \end{bmatrix} \end{bmatrix}, \begin{bmatrix} \begin{bmatrix} 1, \begin{cases} x-47=0\\ y-14=0 \end{bmatrix} \end{bmatrix}, \\ \begin{bmatrix} \begin{bmatrix} 1, \begin{cases} x-47=0\\ y-14=0 \end{bmatrix} \end{bmatrix}, \\ \begin{bmatrix} 14, \begin{cases} x=0\\ y=0 \end{bmatrix} \end{bmatrix} \end{bmatrix}$

The command RegularChains:-TriangularizeWithMultiplicity computes the intersection multiplicities for each point of V(F). In the above Maple session, computations are performed modulo a prime number for the only reason of keeping output expressions small. The same calculations can be performed with the TriangularizeWithMultiplicity command over the reals.

- The theory of regular chains allows us to reduce the question of computing limit points of constructible sets and semi-algebraic sets to that of computing limit points of zero sets of regular chains.
- We will restrict ourselves here to regular chains in dimension 1, that is, where only one variable is free.
- Then, the above question can be solved by computing the Puiseux series solutions of regular chains.