Thesis projects for CS4490

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Research themes and team members

- Symbolic computation: computing exact solutions of algebraic problems on computers with applications to sciences and engineering.
- High-performance computing: making best use of modern computer architectures, in particular hardware accelerators (multi-cores GPUs)

Current students

PhD: Parisa Alvandi, Ning Xie, Mahsa Kazemi, Ruijuan Jing, Xiaohui Chen, Steven Thornton, Robert Moir, Egor Chesakov MSc: Masoud Ataei, Yiming Guan, Davood Mohajerani

Alumni

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Solving polynomial systems symbolically

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Unified (1)+ - Isenser 11 - Maple 14	
o <u>gmait Table D</u> rawing Plot <u>Spreadsheet</u> <u>Tools Window Help</u>	
$R := PolynomialRing([x, y, z]); F := [5*x^2 + 2*x^*z^2 + 5*y^6 + 15*y^4 + 5*z^2 - 15*y^5 polynomial_ring]$	- 5* <i>y</i> ^3];
$5x^{2} + 2xz^{2} + 5y^{6} + 15y^{4} + 5z^{2} - 15y^{5} - 5y^{3}$	m
RealTriangularize(F, R, output = record);	(-/
$\begin{cases} 5 x^2 + 2 z^2 x + 5 y^6 + 15 y^4 - 5 y^3 - 15 y^5 + 5 z^2 = 0\\ 25 y^6 - 75 y^5 + 75 y^4 - z^4 - 25 y^3 + 25 z^2 < 0 \end{cases}$	(2)
$\begin{cases} 5x + z^2 = 0\\ 25y^6 - 75y^5 + 75y^4 - 25y^3 - z^4 + 25z^2 = 0\\ 64z^4 - 1600z^2 + 25 > 0\\ z \neq 0\\ z - 5 \neq 0\\ z + 5 \neq 0 \end{cases}, \begin{cases} x = 0\\ y - 1 = 0\\ z = 0 \end{cases}, \begin{cases} x = 0\\ y = 0\\ z = 0 \end{cases}, \begin{cases} x + 5 = 0\\ y = 0\\ z = 0 \end{cases}, \begin{cases} x = 0\\ y = 0\\ z = 0 \end{cases}, \begin{cases} x = 0\\ y = 0\\ z = 0 \end{cases}, \begin{cases} x = 0\\ y = 0\\ z = 0 \end{cases}$	
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Figure: The *RegularChains* solver designed in our UWO lab is at the heart of Maple, which has about 5,000,000 licences world-wide.

Application to mathematical sciences and engineering

Figure:



Toyota engineers use our software to design control systems

Project 1: Truncated Fourier Transform

- **1** The Fast Fourier Transform (FFT) is a kernel in scientific computing
- 2 It maps a vector of size 2^e to another vector of size 2^e
- The Truncated Fourier Transform (TFT) supports arbitrary vectors but is challenging to implement, in particular on multi/many-cores



FFT with artificial zero points



TFT removes unnecessary computations

Objectives

- Realize an implementation of the TFT and its inverse map
- A configurable Python script will generate the CilkPlus code within the BPAS library www.bpaslib.org

High-performance computing: models of computation



Let \mathbb{K} be the maximum number of thread blocks along an anti-chain of the thread-block DAG representing the program \mathcal{P} . Then the running time $T_{\mathcal{P}}$ of the program \mathcal{P} satisfies:

 $T_{\mathcal{P}} \leq (N(\mathcal{P})/\mathbb{K} + L(\mathcal{P})) C(\mathcal{P}),$

where $C(\mathcal{P})$ is the maximum running time of local operations by a thread among all the thread-blocks, $N(\mathcal{P})$ is the number of thread-blocks and $L(\mathcal{P})$ is the span of \mathcal{P} .

Our UWO lab develops mathematical models to make efficient use of hardware acceleration technology, such as GPUs and multi-core processors. This project is supported by IBM Canada.

Project 2: Models of computation for GPUs

- Several models of computations attempt to estimate the performance of algorithms (or programs) targeting GPGPUs
- The MWP-CWP Model analyzes how computations and memory accesses are interleaved in GPU programs
- The MCM focuses on memory access patterns and memory traffic in GPU algorithms



Objectives

Ompare those models on well-known kernels of scientific computing

2 Can we unify then?

High-performance computing: parallel program translation

```
int main(){
                                                                              void fork_func0(int* sum_a,int* a)
                                       int main()
  int sum_a=0, sum_b=0;
  int a \begin{bmatrix} 5 \end{bmatrix} = \{0, 1, 2, 3, 4\}:
                                                                                      for(int i=0; i<5; i++)</pre>
                                         int sum_a=0, sum_b=0;
  int b[ 5 ] = \{0, 1, 2, 3, 4\};
                                                                                         (*sum a) += a[i]:
                                         int a[5] = \{0,1,2,3,4\};
  #pragma omp parallel
                                         int b[ 5 ] = \{0, 1, 2, 3, 4\};
  ſ
                                                                              void fork func1(int* sum b.int* b)
    #pragma omp sections
                                         meta fork shared(sum a){
                                                                                      for(int i=0; i<5; i++)</pre>
                                           for(int i=0; i<5; i++)</pre>
                                                                                          (*sum_b) += b[ i ];
      #pragma omp section
                                              sum_a += a[ i ];
         for(int i=0: i<5: i++)</pre>
                                                                              int main()
           sum_a += a[ i ];
                                         meta_fork shared(sum_b){
                                                                                int sum_a=0, sum_b=0;
      }
                                           for(int i=0; i<5; i++)</pre>
                                                                                int a[ 5 ] = \{0, 1, 2, 3, 4\};
      #pragma omp section
                                              sum_b += b[ i ];
                                                                                int b[ 5 ] = \{0, 1, 2, 3, 4\};
         for(int i=0: i<5: i++)</pre>
                                                                                cilk_spawn fork_func0(&sum_a,a);
            sum b += b[i]:
                                                                                cilk_spawn fork_func1(&sum_b,b);
                                         meta_join;
      cilk svnc:
}
```

Our lab develops a compilation platform for translating parallel programs from one language to another; above we translate from OpenMP to CilkPlus through MetaFork. This project is supported by IBM Canada.

Project 3: Integrating NPI support into METAFORK

- Currently, the METAFORK language supports different schemes of parallelism: fork-join, pipelining, Single-Instruction Multi-Data.
- CILKPLUS, OPENMP, CUDA code can be generated from METAFORK code by the METAFORK compilation framework





Non-shared memory

Objectives

- Enhance the METAFORK language and METAFORK compilation framework to support non-shared memory and generate MPI code.
- This linguistic extension should be compact while allowing to generate efficient MPI code.

High-performance computing: automatic parallelization

Serial dense univariate polynomial multiplication

```
for(i=0; i<=n; i++){
   for(j=0; j<=n; j++)
        c[i+j] += a[i] * b[j];
}</pre>
```



GPU-like multi-threaded dense univariate polynomial multiplication



We use symbolic computation to automatically translate serial programs to GPU-like programs. This project is supported by IBM Canada.

Project 4: Dependence analysis for parametric GPU kernels

- For performance and portability reasons, GPU kernels should depend on program and machine parameters.
- Standard software tools for automatic parallelization do not support parametric GPU kernels. But METAFORK almost does ...





Input iteration space



Objectives

- Extend the METAFORK framework with a software component for doing dependence analysis on parametric code.
- Note that the METAFORK framework already has the infrastructure to generate parametric GPU kernels.

Research projects with publicly available software

