$\rm UWO\ CS2214$ 

# Tutorial #11

**Problem 1** Let *A* be the set of all ordered pairs of integers for which the second element of the pair is nonzero. Symbolically,

$$A = \mathbf{Z} \times (\mathbf{Z} \setminus \{\mathbf{0}\}).$$

Define a binary relation R on A as follows: For all  $(a, b), (c, d) \in A$ ,

$$(a,b) R (c,d) \Leftrightarrow ac = bd.$$

- 1. Is R reflexive?
- 2. Is R symmetric?
- 3. Is R anti-symmetric?
- 4. Is R transitive?
- 5. Is R an equivalence relation, a partial order, neither, or both?

# Solution 1

1. Is R reflexive? No. Indeed, consider  $(a, c) \in \mathbf{Z} \times (\mathbf{Z} \setminus \{\mathbf{0}\})$ . We have:

$$(a,c) \ R \ (a,c) \Leftrightarrow a^2 = c^2$$

The statement  $a^2 = c^2$  is equivalent to (a - c)(a + c) = 0, that is  $a = c \lor a = -c$ . Therefore, we  $(2,3) \notin R$ . Thus, R is not reflexive.

2. Is R symmetric? Yes. Indeed, consider  $(a, b), (c, d)in \mathbb{Z} \times (\mathbb{Z} \setminus \{0\})$ . We have:

$$(a,b) \mathrel{R} (c,d) \Leftrightarrow ac = bd,$$

and

$$(c,d) R(a,b) \Leftrightarrow ca = db,$$

Clearly, we have:

$$ca = db \Leftrightarrow ac = bd,$$

Therefore, we have:

$$(a,b) R (c,d) \Leftrightarrow (c,d) R (a,b).$$

3. Is R anti-symmetric? No. Indeed, we have (6, 10)R(5, 3).

- 4. Is *R* transitive? No. Indeed, we have (6, 10)R(5, 3) and (5, 3)R(21, 35). But we do **not** have (6, 10)R(21, 35), since  $6 \times 21 \neq 10 \times 35$ .
- 5. Is *R* an equivalence relation, a partial order, neither, or both? Neither. It is not an equivalence relation, since it is not reflexive. It is not a partial order, since it is not anti-symmetric.

**Problem 2** 1. Show that the relation

 $R = \{(x, y) | (x - y) \text{ is an even integer} \}$ 

is an equivalence relation on the set  $\mathbb{R}$  of real numbers.

2. Show that the relation

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R = \{ ((x_1, y_1), (x_2, y_2)) \mid (x_1 < x_2) \text{ or } ((x_1 = x_2) \text{ and } (y_1 \le y_2)) \}
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is a total ordering relation on the set  $\mathbb{R} \times \mathbb{R}$ .

#### Solution 2

- 1. (a) R is <u>reflexive</u>, since for all  $x \in \mathbb{R}$ , we have x x = 0 which is even, hence for all  $x \in \mathbb{R}$ , we have  $(x, x) \in R$ .
  - (b) R is symmetric, since for all  $x, y \in \mathbb{R}$ , if  $x y \equiv 0 \mod 2$  holds then so does  $y - x \equiv 0 \mod 2$ , that is, if  $(x, y) \in R$  holds then so does  $(y, x) \in R$ .
  - (c) R is <u>transitive</u>, since for all  $x, y, z \in \mathbb{R}$ , if  $x y \equiv 0 \mod 2$  and  $y x \equiv 0 \mod 2$  both hold then so does  $x z = (x y) + (y z) \equiv 0 \mod 2$ , that is, if  $(x, y) \in R$  and  $(y, z) \in R$  both hold then so does  $(x, z) \in R$ .

Therefore, R is an equivalence relation.

- 2. (a) R is <u>reflexive</u>, since for all  $(x_1, y_1) \in \mathbb{R} \times \mathbb{R}$ , we have  $((x_1 = x_1) \text{ and } y_1 \leq y_1$ , that is, for all  $(x_1, y_1) \in \mathbb{R} \times \mathbb{R}$  we have  $((x_1, y_1), (x_1, y_1)) \in \mathbb{R}$ .
  - (b) R is anti-symmetric, since for all  $(x_1, y_1), (x_2, y_2) \in \mathbb{R} \times \mathbb{R}$ , if  $((x_1, y_1), (x_2, y_2)) \in R$  and  $((x_2, y_2), (x_1, y_1)) \in R$  both hold then neither  $x_1 < x_2$  nor  $x_2 < x_1$  holds but both  $((x_1 = x_2) \text{ and } y_1 \leq y_2)$  and  $((x_2 = x_1) \text{ and } y_2 \leq y_1)$  hold, which implies  $(x_1, y_1) = (x_2, y_2)$ .

(c) R is <u>transitive</u>. To prove this consider  $(x_1, y_1), (x_2, y_2), (x_3, y_3) \in \mathbb{R} \times \mathbb{R}$  such that  $((x_1, y_1), (x_2, y_2)) \in R$  and  $((x_2, y_2), (x_3, y_3)) \in R$  both hold. We shall prove that  $((x_1, y_1), (x_3, y_3)) \in R$  also holds. Four cases must be inspected:

i.  $x_1 < x_2$  and  $x_2 < x_3$ , ii.  $x_1 < x_2$  and  $x_2 = x_3$  and  $y_2 \le y_3$ , iii.  $x_1 = x_2$  and  $y_1 \le y_2$  and  $x_2 < x_3$ , iv.  $x_1 = x_2$  and  $y_1 \le y_2$  and  $x_2 = x_3$  and  $y_2 \le y_3$ , which respectively imply:

i.  $x_1 < x_3$ , ii.  $x_1 < x_3$ , iii.  $x_1 < x_3$ , iv.  $x_1 = x_3$  and  $y_1 \le y_3$ , that is  $((x_1, y_1), (x_3, y_3)) \in R$ .

- 3. Therefore, R is an ordering relation on the set  $\mathbb{R} \times \mathbb{R}$ .
- 4. *R* is a <u>total</u> ordering relation on the set  $\mathbb{R} \times \mathbb{R}$ . Indeed, for all  $(x_1, y_1), (x_2, y_2) \in \mathbb{R} \times \mathbb{R}$ , we have
  - (a) either  $x_1 < x_2$  (in which case  $((x_1, y_1), (x_2, y_2)) \in R$  holds),
  - (b) or  $(x_1 = x_2 \text{ and } y_1 \leq y_2 \text{ (in which case } ((x_1, y_1), (x_2, y_2)) \in R \text{ holds}),$
  - (c) or  $(x_1 = x_2 \text{ and } y_1 > y_2$  (in which case  $((x_2, y_2), (x_1, y_1)) \in R$  holds),
  - (d) or  $x_1 > x_2$  (in which case  $((x_2, y_2), (x_1, y_1)) \in R$  holds).

**Problem 3** Let R be a binary relation on a set A. We denote by I the *identity relation* on A, that is:

$$I = \{ (x, x) \mid x \in A \}.$$

We denote by r(R) the relation given by:

$$r(R) = R \cup I.$$

- 1. Prove that r(R) is reflexive.
- 2. Prove that R is reflexive if and only if r(R) = R.

Clearly, if R' is a reflexive relation so that  $R \subseteq R'$  holds then  $r(R) \subseteq R'$  holds as well. For that reason, the relation r(R) can be regarded as the "smallest" reflexive relation containing R and r(R) is called the *reflexive closure* of R.

## Solution 3

- 1. Indeed R reflexive exactly means  $I \subseteq R$ .
- 2. From the previous question, if R reflexive, then  $I \subseteq R$  holds and thus  $r(R) \subseteq R$  holds. Since  $R \subseteq r(R)$  clearly holds as well, we have proved the following:

$$R$$
 reflexive  $\rightarrow r(R) = R$ 

The converse follow from the previous question.

**Problem 4** Let R be a binary relation on a set A. We denote by  $R^{-1}$  the *inverse relation* of R, that is, the binary relation on A defined by:

$$R^{-1} = \{ (y, x) \mid (x, y) \in R \}.$$

We denote by s(R) the relation given by:

$$s(R) = R \cup R^{-1}.$$

- 1. Prove that s(R) is symmetric.
- 2. Prove that R is symmetric if and only if s(R) = R.
- 3. Prove that if R' is a symmetric relation so that  $R \subseteq R'$  holds, then  $s(R) \subseteq R'$  holds as well.

From the third question it follows that the relation s(R) can be regarded as the "smallest" symmetric relation containing R. For that reason, s(R) is called the symmetric closure of R.

#### Solution 4

1. Let us prove that s(R) is symmetric, thus let us prove that for all  $x, y \in A$ , if  $(x, y) \in s(R)$ , then  $(y, x) \in s(R)$  holds as well. So, let  $x, y \in A$  and assume that  $(x, y) \in s(R)$  holds. Since  $s(R) = R \cup R^{-1}$  holds, two cases arise: either  $(x, y) \in R$  holds or  $(x, y) \in R^{-1}$  holds. Consider the first case. Then, by definition of  $R^{-1}$ , we have  $(y, x) \in R^{-1}$ , thus we have  $(y, x) \in s(R)$ . Consider now the second case, that is,  $(x, y) \in R^{-1}$ . Then, by definition of  $R^{-1}$ , we have  $(y, x) \in R$ , thus we have  $(y, x) \in s(R)$ . Finally, we have shown that s(R) is symmetric.

- 2. Let us prove that R is symmetric if and only if s(R) = R. First, we assume that R is symmetric and we prove that s(R) = R holds as well. We observe that R symmetric implies that  $R^{-1} \subseteq R$  holds and thus we have s(R) = R. Conversely, if s(R) = R holds, then  $R^{-1} \subseteq R$  holds as well which implies that R is symmetric.
- 3. Let R' be a symmetric relation so that  $R \subseteq R'$  holds. We shall prove that  $s(R) \subseteq R'$  holds as well. Since  $R \subseteq R'$  holds, it is a routine exercise to prove that  $s(R) \subseteq s(R')$  holds as well. Since R' is symmetric, it follows from the second question that R' = s(R'). Therefore, we have  $s(R) \subseteq R'$ , as required.

**Problem 5** Let R be a binary relation on a finite set A with cardinality n. We denote by t(R) the *transitive closure* of R, that is, the binary relation on A defined by:

$$t(R) = R \cup R^2 \cup \dots \cup R^n.$$

1. Let k be an integer such that  $2 \le k \le n$ . Let x, y be in A. We denote by P(x, y, k) the following predicate:

there exist (k-1) elements  $x_2, \ldots, x_k$  of A so that  $(x, x_2), (x_2, x_3), \ldots, (x_k, y)$  all belong to R.

Prove that the following statements are equivalent for all  $x, y \in A$ :

- (a)  $(x, y) \in \mathbb{R}^k$ ,
- (b) P(x, y, k) holds
- 2. Let  $k, \ell$  be two positive integers, with  $k \leq n$  and  $\ell \leq n$ . Let x, y, z be in A so that P(x, y, k) and  $P(y, z, \ell)$  both hold. Prove that P(x, z, m), with  $m = \min(n, k + \ell)$ , also holds.
- 3. Prove that t(R) is transitive.
- 4. Prove that if R transitive, then  $R^k \subseteq R$  for all positive integer k.
- 5. Prove that R transitive if and only if t(R) = R.
- 6. Prove that if R' is a transitive relation so that  $R \subseteq R'$  holds, then  $t(R) \subseteq R'$  holds as well.

It follows from the last question that the relation t(R) can be regarded as the "smallest" transitive relation containing R. This is the reason why t(R)is called the *transitive closure* of R.

## Solution 5

- 1. We proceed by induction on k, for  $1 \leq k \leq n$ . We observe that the equivalence (a)  $\iff$  (b) is clear for all  $x, y \in A$ , when k = 1. Indeed, P(x, y, 1) simply means  $(x, y) \in R$ . Now we assume that for some k, with  $1 \leq k < n$ , the equivalence (a)  $\iff$  (b) holds for all  $x, y \in A$ . We shall prove that this equivalence holds for all  $x, y \in A$ , with k + 1 instead of k. So let  $x, y \in A$ . Assume first that  $(x, y) \in \mathbb{R}^{k+1}$  holds and let us prove that P(x, y, k+1) holds as well. By definition of  $R^{k+1}$ , we have  $R^{k+1} = R \circ R^k$ , thus there exists  $z \in A$ so that  $(x, z) \in \mathbb{R}^k$  and  $(z, y) \in \mathbb{R}$ . By induction hypothesis,  $(x, z) \in \mathbb{R}^k$  $R^k$  is equivalent to P(x, z, k), that is, there exist (k - 1) elements  $x_2, \ldots, x_k$  of A so that  $(x, x_2), (x_2, x_3), \ldots, (x_k, z)$  all belong to R. Putting everything together, we deduce that there exist k elements  $x_2, \ldots, x_k, z$  of A so that  $(x, x_2), (x_2, x_3), \ldots, (x_k, z), (z, y)$  all belong to R. This latter statement means that P(x, y, k+1) holds, as required. Proving the converse implication (that is,  $P(x, y, k+1) \rightarrow (x, y) \in$  $R^{k+1}$ ) can easily be done using the same arguments as those used for proving the direct implication  $(x, y) \in \mathbb{R}^{k+1} \to \mathbb{P}(x, y, k+1)$ . This completes the proof of this first question.
- 2. Let  $k, \ell$  be two positive integers, with  $k \leq n$  and  $\ell \leq n$ . Let x, y, z be in A so that P(x, y, k) and  $P(y, z, \ell)$  both hold. We shall prove that P(x, z, m), with  $m = \min(n, k + \ell)$ , holds as well. Recall first that P(x, y, k) means that there exist (k 1) elements  $x_2, \ldots, x_k$  of A so that  $(x, x_2), (x_2, x_3), \ldots, (x_k, y)$  all belong to R. Similarly,  $P(y, z, \ell)$  means that there exist  $(\ell 1)$  elements  $x_{k+2}, \ldots, x_{\ell+k}$  so that  $(y, x_{k+2}), \ldots, (x_{\ell+k}, z)$  all belong to R. It follows that there exist  $x_2, \ldots, x_k, x_{k+1}, x_{k+2}, \ldots, x_{\ell+k} \in A$  with  $y = x_{k+1}$ , so that

$$(x, x_2), (x_2, x_3), \dots, (x_k, x_{k+1}), (x_{k+1}, x_{k+2}), \dots, (x_{\ell+k}, z)$$

all belong to R. The number of these "intermediate points"

$$x_2,\ldots,x_k,x_{k+1},x_{k+2},\ldots,x_{\ell+k}$$

is  $\ell - k - 1$ . But if  $\ell - k - 1$  exceeds n - 1 then there is necessarily some repetitions among those points and thus some arcs can be removed.

Indeed, since the set A counts n elements, the number of these "intermediate points" (excluding x and z) is at most n-1 if x = z holds and n-2 otherwise. Therefore, the number of intermediate points is m-1with  $m = \min(n, k + \ell)$ . Therefore, we have P(x, z, m), as required.

3. Let us prove that t(R) is transitive. Let x, y, z be in A so that  $(x, y) \in t(R)$  and  $(y, z) \in t(R)$  both hold. Let us prove that  $(x, z) \in t(R)$  as well. Recall that, by definition of t(R), we have:

$$t(R) = R \cup R^2 \cup \dots \cup R^n.$$

Therefore, the statement  $(x, y) \in t(R)$  means that there exists a positive integer  $k \leq n$  so that  $(x, y) \in R^k$ . Similarly, the statement  $(y, z) \in t(R)$  means that there exists a positive integer  $\ell \leq n$  so that  $(y, z) \in R^{\ell}$ . From the first question, we deduce that P(x, y, k) and  $P(y, z, \ell)$  both hold. Then, from the second question, we deduce that P(x, z, m), with  $m = \min(n, k + \ell)$ , also holds. This implies, using the first question again that  $(x, z) \in R^m$ . Since  $m \leq n$  holds, it follows that (x, z) belongs to one of  $R, R^2, \ldots, R^n$ . In other words, (x, z)belongs to t(R), as required. This completes the proof that t(R) is transitive.

4. The proof is by induction  $k \ge 1$ . The base step k = 1 is clear since we obviously have  $R \subseteq R$ . We now prove the *inductive step*. We assume that  $R^k \subseteq R$  holds for some  $k \ge 1$ . We shall prove that  $R^{k+1} \subseteq R$  holds as well. Recall that we have  $R^{k+1} = R \circ R^k$ . Since  $R^k \subseteq R$  holds (by induction hypothesis) a routine proof yields

$$R \circ R^k \subseteq R \circ R.$$

Since R is transitive, it follows directly from the definition of the composition of two relations that  $R \circ R \subseteq R$  holds. Therefore, we have  $R^{k+1} \subseteq R$ , which completes the proof of the inductive step and thus the proof of the fact that if R transitive, then  $R^k \subseteq R$  for all positive integer k.

5. We prove the equivalence:

$$R$$
 transitive  $\iff t(R) = R.$ 

We first assume that R is transitive. Recall that we have:

$$t(R) = R \cup R^2 \cup \dots \cup R^n.$$

From the previous question, we have  $R^k \subseteq R$ , for all positive integer  $k \geq 1$ . This clearly implies t(R) = R. Conversely, if t(R) = R holds, then from the third question, we deduce that R is transitive, as required.

6. Let R' be a transitive relation so that  $R \subseteq R'$  holds. We prove that  $t(R) \subseteq R'$  holds as well. From  $R \subseteq R'$ , an easy routine proof (similar to the proof of the fourth question) yields  $t(R) \subseteq t(R')$ . Since R' is transitive, the fifth question yields t(R') = R'. Therefore, we have  $t(R) \subseteq R'$ , as required.